Data Mining Techniques

CS 6220 - Section 3 - Fall 2016

Lecture 9

Jan-Willem van de Meent (*credit*: Yijun Zhao, Carla Brodley, Eamon Keogh, Tan + Steinbach + Kumar)



Hierarchical

Partitional





Create a hierarchical decomposition using *"some criterion"*

Construct partitions and evaluate them using *"some criterion"*

1. Connectivity-based (Hierarchical)



Notion of Clusters: Cut off dendrogram at some depth

2. Centroid-based (K-means, K-medoids)



Notion of Clusters: Voronoi tesselation

3. Density-based (DBSCAN, OPTICS)



Notion of Clusters: Connected regions of high density

4. Distribution-based (Mixture Models)



Notion of Clusters: Distributions on features



Idea: Minimize Sum of Squares

$$SSE_i = \sum_{\mathbf{x} \in C_i} \| \mathbf{x} - \mu_i \|^2$$

 $SSE = \sum_{j=1}^{K} SSE_j$

Brute force search: NP-hard How many partitions?



Idea: Minimize Sum of Squares

$$SSE_i = \sum_{\mathbf{x} \in C_i} \| \mathbf{x} - \mu_i \|^2$$

 $SSE = \sum_{j=1}^{K} SSE_j$

Use *heuristic* search (as in hierarchical case)



Randomly initialize K centroids μ_k



Assign each point to closest centroid, then update centroids to average of points



Assign each point to closest centroid, then update centroids to average of points



Repeat until convergence (no points reassigned, means unchanged)



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K-means Algorithm

Input:

 $X = \{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_N\}$ Number of clusters *K*

Initialize: K random centroids $\mu_1, \mu_2, \ldots, \mu_K$ Repeat Until Convergence

1 For
$$i = 1, ..., K$$
 do
 $C_i = \{\mathbf{x} \in X | i = \arg \min_{1 \le j \le K} \| \mathbf{x} - \mu_j \|^2\}$
2 For $i = 1, ..., K$ do
 $\mu_i = \arg \min_{\mathbf{z}} \sum_{\mathbf{x} \in C_i} \| \mathbf{z} - \mathbf{x} \|^2$
Output: $C_1, C_2, ..., C_K$

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Output: $C_1, C_2, ..., C_K$

- K-means: Set μ to mean of points in C
- K-medoids: Set $\mu = x$ for point in C with minimum SSE

K-means Complexity

Cost function

$$Cost = \sum_{1}^{K} \sum_{\mathbf{x} \in \mathbf{C}_{i}} \| \mathbf{x} - \boldsymbol{\mu}_{i} \|^{2}$$

What is the computational complexity (per iteration) for *N* points with *D features* and *K* clusters?

K-means Complexity

Cost function

$$Cost = \sum_{1}^{K} \sum_{\mathbf{x} \in \mathbf{C}_{i}} \| \mathbf{x} - \boldsymbol{\mu}_{i} \|^{2}$$

Complexity: O(N K D #iterations)

Web-scale k-means clustering

D Sculley - Proceedings of the 19th international conference on ..., 2010 - dl.acm.org Abstract We present two modifications to the popular k-means clustering algorithm to address the extreme requirements for latency, scalability, and sparsity encountered in userfacing web applications. First, we propose the use of mini-batch optimization for k-means ... Cited by 152 Related articles All 11 versions Cite Save

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Web-Scale K-Means Clustering

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ABSTRACT

We present two modifications to the popular k-means clustering algorithm to address the extreme requirements for latency, scalability, and sparsity encountered in user-facing web applications. First, we propose the use of mini-batch optimization for k-means clustering. This reduces computation cost by orders of magnitude compared to the classic batch algorithm while vielding significantly better solutions than online stochastic gradient descent. Second, we achieve sparsity with projected gradient descent, and give a fast ϵ accurate projection onto the L1-ball. Source code is freely available: http://code.google.com/p/sofia-ml

Categories and Subject Descriptors

I.5.3 [Computing Methodologies]: Pattern Recognition-Clustering

General Terms

Algorithms, Performance, Experimentation

Keywords

unsupervised clustering, scalability, sparse solutions

1. CLUSTERING AND THE WEB

Unsupervised clustering is an important task in a range of web-based applications, including grouping search results, near-duplicate detection, and news aggregation to name but a few. Lloyd's classic k-means algorithm remains a popular choice for real-world clustering tasks [6]. However, the standard batch algorithm is slow for large data sets. Even optimized batch k-means variants exploiting triangle inequality [3] cannot cheaply meet the latency needs of user-facing applications when clustering results on large data sets are required in a fraction of a second.

This paper proposes a mini-batch k-means variant that yields excellent clustering results with low computation cost on large data sets. We also give methods for learning sparse cluster centers that reduce storage and network cost.

2. MINI-BATCH K-MEANS

The k-means optimization problem is to find the set C of cluster centers $\mathbf{c} \in \mathbb{R}^m$, with |C| = k, to minimize over a set

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WWW 2010, April 26–30, 2010, Raleigh, North Carolina, USA. ACM 978-1-60558-799-8/10/04.

X of examples $\mathbf{x} \in \mathbb{R}^m$ the following objective function:

 $\min \sum_{\mathbf{x} \in X} ||f(C, \mathbf{x}) - \mathbf{x}||^2$

Here, $f(C, \mathbf{x})$ returns the nearest cluster center $\mathbf{c} \in C$ to \mathbf{x} using Euclidean distance. It is well known that although this problem is NP-hard in general, gradient descent methods converge to a local optimum when seeded with an initial set of k examples drawn uniformly at random from X [1].

The classic batch k-means algorithm is expensive for large data sets, requiring O(kns) computation time where n is the number of examples and s is the maximum number of nonzero elements in any example vector. Bottou and Bengio proposed an online, stochastic gradient descent (SGD) variant that computed a gradient descent step on one example at a time [1]. While SGD converges quickly on large data sets, it finds lower quality solutions than the batch algorithm due to stochastic noise [1].

Algorithm 1 Mini-batch k-Means.

1:	Given: k , mini-batch size b , iterations t , data set X
2:	Initialize each $\mathbf{c} \in C$ with an \mathbf{x} picked randomly from X
3:	$\mathbf{v} \leftarrow 0$
4:	for $i = 1$ to t do
5:	$M \leftarrow b$ examples picked randomly from X
6:	for $x \in M$ do
7:	$\mathbf{d}[\mathbf{x}] \leftarrow f(C, \mathbf{x})$ // Cache the center nearest to \mathbf{x}
8:	end for

for $\mathbf{x} \in M$ do $\mathbf{c} \gets \mathbf{d}[\mathbf{x}]$ // Get cached center for this x $\mathbf{v}[\mathbf{c}] \leftarrow \mathbf{v}[\mathbf{c}] + 1$ // Update per-center counts // Get per-center learning rate $\eta \leftarrow \frac{1}{\mathbf{v}[\mathbf{c}]}$ $\mathbf{c} \leftarrow (\mathbf{i} - \eta)\mathbf{c} + \eta \mathbf{x}$ // Take gradient step 14: end for 15: end for

We propose the use of mini-batch optimization for k-means clustering, given in Algorithm 1. The motivation behind this method is that mini-batches tend to have lower stochastic noise than individual examples in SGD (allowing convergence to better solutions) but do not suffer increased computational cost when data sets grow large with redundant examples. We use per-center learning rates for fast convergence, in the manner of [1]; convergence properties follow closely from this prior result [1].





Figure 1: Convergence Speed. The mini-batch method (blue) is orders of magnitude faster than the full batch method (green), while converging to significantly better solutions than the online SGD method (red).

our experiments. To assess performance at scale, the set of 781,265 examples were used for training and the remaining 23,149 examples for testing. On each trial, the same random initial cluster centers were used for each method. We evaluated the learned cluster centers using the k-means objective function on the held-out test set; we report fractional error from the best value found by the batch algorithm run. to convergence. We set the mini-batch b to 1000 based on separate initial tests; results were robust for a range of b.

The results (Fig. 1) show a clear win for mini-batch kmeans. The mini-batch method converged to a near optimal value several orders of magnitude faster than the full batch method, and also achieved significantly better solutions than SGD. Additional experiments (omitted for space) showed that mini-batch k-means is several times faster on large data sets than batch k-means exploiting triangle inequality [3].

For small values of k, the mini-batch methods were able to produce near-best cluster centers for nearly a million documents in a fraction of a CPU second on a single ordinary 2.4 GHz machine. This makes real-time clustering practical for user-facing applications.

3. SPARSE CLUSTER CENTERS

We modify mini-batch k-means to find sparse cluster centers, allowing for compact storage and low network cost. The intuition for seeking sparse cluster centers for document clusters is that term frequencies follow a power-law distribution. Many terms in a given cluster will only occur in one or two documents, giving them very low weight in the cluster center. It is likely that for a locally optimal center \mathbf{c} , there is a nerby point \mathbf{c}' with many fewer non-zero values.

Sparsity may be induced in gradient descent using the projected-gradient method, projecting a given \mathbf{v} to the nearest point in an L1-ball of radius λ after each update [2]. Thus, for mini-batch k-means we achieve sparsity by performing an L1-ball projection on each cluster center **c** after each mini-batch iteration.

Algorithm 2 ϵ -L1: an ϵ -Accurate Projection to L1 Ball. 1: Given: ϵ tolerance, L1-ball radius λ , vector $\mathbf{c} \in \mathbb{R}^n$

2: if $||\mathbf{c}||_i \leq \lambda + \epsilon$ then exit

- 3: upper $\leftarrow ||\mathbf{c}||_{\infty}$; lower $\leftarrow 0$; current $\leftarrow ||\mathbf{c}||_1$
- 4: while $current > \lambda(1 + \epsilon)$ or $current < \lambda$ do $\theta \leftarrow \frac{upper+lower}{20}$ // Get L1 value for this θ 5:
- $\begin{array}{l} \theta \leftarrow \frac{p_{I}(\mathbf{r})}{2.0} & // \text{ Get } L1 \text{ value for th} \\ current \leftarrow \sum_{\mathbf{c}_i \neq 0} \max(0, |\mathbf{c}_i| \theta) \\ \text{if } current \leq \lambda \text{ then } upper \leftarrow \theta \text{ else } lower \leftarrow \theta \end{array}$ 6:
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- 8. end while
- 9: for *i* = 1 to *m* do
- 10: $\mathbf{c}_i \leftarrow sign(\mathbf{c}_i) \ast \max(0, |\mathbf{c}_i| \theta) // \text{ Do the projection}$ 11: end for

Fast L1 Projections. Applying L1 constraints to kmeans clustering has been studied in forthcoming work by Witten and Tibshirani [5]. There, a hard L1 constraint was applied in the full batch setting of maximizing betweencluster distance for k-means (rather than minimizing the k-means objective function directly); the work did not discuss how to perform this projection efficiently.

The projection to the L1 ball can be performed effectively using, for example, the linear time L1-ball projection algorithm of Duchi et al. [2], referred to here as LTL1P. We give an alternate method in Algorithm 2, observing that the exact L1 radius is not critical for sparsity. This simple approximation algorithm uses bisection to find a value θ that projects **c** to an L1 ball with radius between λ and $(1 + \epsilon)\lambda$. Our method is easy to implement and is also significantly faster in practice than LTL1P due to memory concurrency.

METHOD	λ	#NON-ZERO'S	TEST OBJECTIVE	CPUs
full batch	-	200,319	0 (baseline)	133.96
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Results. Using the same set-up as above, we tested Duchi et al.'s linear time algorithm and our ϵ -accurate projection for mini-batch k-means, with a range of λ values. The value of ϵ was arbitrarily set to 0.01. We report values for k = 10, b = 1000, and t = 16 (results for other values qualitatively similar). Compared with the full batch method, we achieve much sparser solutions. The approximate projection is roughly twice as fast as LTL1P and finds sparser solutions. but gives slightly worse performance on the test set. These results show that sparse clustering may cheaply be achieved with low latency for user-facing applications.

4. REFERENCES

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Experiments. We tested the mini-batch k-means against both Lloyd's batch k-means [6] and the SGD variant of [1]. We used the standard RCV1 collection of documents [4] for



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Algorithm 1 Mini-batch k-Means.

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4: for
$$i = 1$$
 to t do

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$$\mathbf{x} \in M$$
 do

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9: for
$$\mathbf{x} \in M$$
 do

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13:
$$\mathbf{c} \leftarrow (1 - \eta)\mathbf{c} + \eta \mathbf{x}$$
 // Take gradient step

```
end for
14:
```

```
15: end for
```

Complexity: O(H M K D t)





"Good" initial choice



"Bad" initial choice

What is the chance of randomly selecting one point from each of *K* clusters?

(assume each cluster has size n = N/K)

What is the chance of randomly selecting one point from each of *K* clusters?

(assume each cluster has size n = N/K)

 $\frac{\text{ways to select one from each cluster}}{\text{ways to select K centroids}} = \frac{K!n^{K}}{(Kn)^{K}} = \frac{K!}{K^{K}}$

Implication: We will almost always have multiple initial centroids in same cluster.



5 pairs of clusters, two initial points in each pair



Some pairs have 3 initial centroids, some have 1



Conclusion: Heuristic search leads to local optima

Initialization tricks

- Use multiple restarts
- Initialize with hierarchical clustering
- Select more than K points, keep most widely separated points

Choosing K



K=3, SSE=134



Choosing K



"Elbow finding" (a.k.a. "knee finding") Set K to value just above "abrupt" increase

K-means Limitations: Differing Sizes



K-means Limitations: Differing Sizes



K-means Limitations: Differing Densities


K-means Limitations: Differing Densities



K-means Limitations: Non-globular Shapes



K-means Limitations: Non-globular Shapes



Overcoming K-means Limitations



Intuition: "Combine" smaller clusters into larger clusters

- One Solution: Hierarchical Clustering
- Another Solution: Density-based Clustering

Density-based Clustering

DBSCAN



arbitrarily shaped clusters

[PDF] A density-based algorithm for discovering clusters in large spatial databases with noise. <u>M Ester, HP Kriegel, J Sander, X Xu</u> - Kdd, 1996 - aaai.org Abstract Clustering algorithms are attractive for the task of class identification in spatial databases. However, the application to large spatial databases rises the following requirements for clustering algorithms: minimal requirements of domain knowledge to ... Cited by 8901 Related articles All 70 versions Cite Save More

(one of the most-cited clustering methods)

DBSCAN



arbitrarily shaped clusters

Intuition

- A *cluster* is a region of *high* density
- Noise points lie in regions of low density

Naïve approach

For each point in a cluster there are at least a minimum number (MinPts) of points in an Eps-neighborhood of that point.



cluster

Eps-neighborhood of a point p

 $N_{Eps}(p) = \{ q \in D \mid dist (p, q) \le Eps \}$



- In each cluster there are two kinds of points:
 - points inside the cluster (core points)
 - points on the border (border points)



cluster

An Eps-neighborhood of a border point contains significantly less points than an Eps-neighborhood of a core point.

Better notion of cluster

For every point p in a cluster C there is a point $q \in C$, so that

(1) p is inside of the Eps-neighborhood of q

and

(2) $N_{Eps}(q)$ contains at least MinPts points.



Density Reachability

Definition

A point p is directly density-reachable from a point q with regard to the parameters Eps and MinPts, if

1) $p \in N_{Eps}(q)$ (reachability)

2)

 $|N_{Eps}(q)| \ge MinPts$ (core point condition)



Parameter: MinPts = 5 p directly density reachable from q $p \in N_{Eps}(q)$ $|N_{Eps}(q)| = 6 \ge 5 = MinPts$ (core point condition) q not directly density reachable from p $|N_{Eps}(p)| = 4 < 5 = MinPts$ (core point condition)

Note: This is an asymmetric relationship

Density Reachability

Definition

A point p is density-reachable from a point q with regard to the parameters Eps and MinPts if there is a chain of points $p_1, p_2, ..., p_s$ with $p_1 = q$ and $p_s = p$ such that p_{i+1} is directly density-reachable from p_i for all 1 < i < s-1.



MinPts = 5

 $|N_{Eps}(q)| = 5 = MinPts$ (core point condition)

 $|N_{Eps}(p_1)| = 6 \ge 5 = MinPts$ (core point condition)

Density Connectivity

Definition (density-connected)

A point p is density-connected to a point q

with regard to the parameters Eps and MinPts

if there is a point v such that both p and q are density-reachable from v.



Note: This is a symmetric relationship

Definition of a Cluster

A cluster with regard to the parameters Eps and MinPts is a non-empty subset C of the database D with



2) For all p, $q \in C$:

(Connectivity)

The point p is density-connected to q with regard to the parameters Eps and MinPts.

Definition of Noise

Let $C_1,...,C_k$ be the clusters of the database D with regard to the parameters Eps_i and MinPts₁ (i=1,...,k).

The set of points in the database D not belonging to any cluster C_1, \dots, C_k is called **noise**:

Noise = { $p \in D | p \notin C_i$ for all i = 1,...,k }



DBSCAN Algorithm

 Start with an arbitrary point p from the database and retrieve all points density-reachable from p with regard to Eps and MinPts.

(2) If p is a core point, the procedure yields a cluster with regard to Eps and MinPts and all points in the cluster are classified.

(3) If p is a **border point**, no points are density-reachable from p and DBSCAN visits the next unclassified point in the database.

DBSCAN Algorithm



Original Points

Point types: core, border and noise

DBSCAN Complexity

- Time complexity: O(N²) if done naively, O(N log N) when using a spatial index (works in relatively low dimensions)
- Space complexity: O(N)

DBSCAN strengths



- + Resistant to noise
- + Can handle arbitrary shapes

DBSCAN Weaknesses



Ground Truth

MinPts = 4, Eps=9.92 *Min*

MinPts = 4, Eps=9.75

- Varying densities
- High dimensional data
- Overlapping clusters

Determining EPS and MINPTS



- Calculate distance of k-th nearest neighbor for each point
- Plot in ascending / descending order
- Set EPS to max distance before "jump"

K-means vs DBSCAN



Evaluation of Clustering

Clusters in Random Data



Clustering Criteria

• External Quality Criteria

- Precision-Recall Measure
- Mutual Information

- Internal Quality Criteria
 Measure compactness of clusters
 - Sum of Squared Error (SSE)
 - Scatter Criteria

 $I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$

$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

Uncorrelated Variables

p(a,b) = p(a)p(b)

$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

Uncorrelated Variables

p(a,b) = p(a)p(b)

$$I(A;B) = \sum_{a \in A, b \in B} p(a)p(b)\log \frac{p(a)p(b)}{p(a)p(b)} = 0$$

$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

$$p(A = a, B = b) = \delta(a, b)p(B = b)$$

$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

Perfectly Correlated Variables

$$p(A = a, B = b) = \delta(a, b)p(B = b)$$

 $I(A;B) = \sum_{a \in A, B \in B} p(A = a, B = b) \log \frac{p(A = a, B = b)}{p(A = a)p(B = b)}$

$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

$$p(A = a, B = b) = \delta(a, b)p(B = b)$$

$$I(A;B) = \sum_{b \in B} p(B=b) \log \frac{p(B=b)}{p(A=b)p(B=b)}$$

$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

$$P(A=k) = \sum_{l \in B} \delta(k,l)p(B=l) = p(B=k)$$

$$I(A;B) = \sum_{b \in B} p(B=b) \log \frac{p(B=b)}{p(A=b)p(B=b)}$$

$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

$$P(A=k) = \sum_{l \in B} \delta(k,l)p(B=l) = p(B=k)$$

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$$I(A,B) = \sum_{a \in A, b \in B} p(a,b) \log \frac{p(a,b)}{p(a)p(b)}$$

$$P(A=k) = \sum_{l \in B} \delta(k,l)p(B=l) = p(B=k)$$

$$I(A;B) = -\sum_{b \in B} p(b) \log p(b) = H(B)$$

 $I(Y;Z) = \sum_{y,z} p(y,z) \log \frac{p(y,z)}{p(y)p(z)}$

y_n: True class label for example *n z_n*: Clustering label for example *n*
$$I(Y;Z) = \sum_{y,z} p(y,z) \log \frac{p(y,z)}{p(y)p(z)}$$

y_n: True class label for example *n z_n*: Clustering label for example *n*

$$p(Y = k) = \frac{1}{N} \sum_{n} I(y_n = k)$$
 $p(Z = l) = \frac{1}{N} \sum_{n} I(z_n = l)$

$$p(Y = k, Z = l) = \frac{1}{N} \sum_{n} I(y_n = k \land z_n = l)$$

 $I(Y;Z) = \sum_{y,z} p(y,z) \log \frac{p(y,z)}{p(y)p(z)}$

p(y,z)	1	2	3	<i>p</i> (<i>y</i>)
cat	0.39	0.08	0.02	0.49
dog	0.06	0.31	0.01	0.38
parrot	0.01	0.01	0.11	0.13
p(z)	0.46	0.40	0.04	





What happens to I(Y;Z) if we swap cluster labels?

 $I(Y;Z) = \sum_{y,z} p(y,z) \log \frac{p(y,z)}{p(y)p(z)}$

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p(z)	0.40	0.46	0.04	

Mutual Information is *invariant* under label permutations

Scatter Criteria (Internal)

Let

Define

- $\mathbf{x} = (x_1, \dots, x_d)^T$ $C_1, \dots, C_K \text{ be a clustering of } \{\mathbf{x}_1, \dots, \mathbf{x}_N\}$
- Size of each cluster:

$$N_i = |C_i| \qquad i = 1, 2, \ldots, K$$

• Mean for each cluster:

$$\mu_i = \frac{1}{N_i} \sum_{x \in C_i} \mathbf{x} \qquad i = 1, 2, \dots, K$$

• Total mean : $\mu = \frac{1}{N} \sum_{i=1}^{N} \mathbf{x}_{i} \quad \text{OR} \quad \mu = \frac{1}{N} \sum_{i=1}^{K} N_{i} \mu_{i}$

Scatter Criteria (Internal)

• Scatter matrix for the *i*th cluster:

$$S_i = \sum_{\mathbf{x} \in C_i} (\mathbf{x} - \mu_i) (\mathbf{x} - \mu_i)^T$$
 (outer product)

- Within cluster scatter matrix : $S_W = \sum_{i=1}^{K} S_i$
- Between cluster scatter matrix : $S_B = \sum_{i=1}^{K} N_i (\mu_i - \mu) (\mu_i - \mu)^T$ (outer product)

Scatter Criteria (Internal)

- The trace criteria: sum of the diagonal elements of a matrix
- A good partition of the data should have:
 Low tr(S_W): similar to minimizing SSE
 - High $tr(S_B)$

• High
$$\frac{tr(S_B)}{tr(S_W)}$$