# L05: SQL

### Announcements!

- HW1 is due tonight
- HW2 groups are assigned
- Outline today:
  - nested queries and witnesses
  - We start with a detailed example!
  - outer joins, nulls?

# Small IMDB schema (SQLite)





# Big IMDB schema (Postgres)



# Theta joins







A **Theta-join** allows for arbitrary comparison relationships (such as  $\geq$ ). An **equijoin** is a theta join using the equality operator.

# Theta joins



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Witnesses: with joins (1/6)





Q: <u>For each company</u>, find the most expensive product + its price

Witnesses: with joins (2/6)



Q: For each company, find the most expensive product + its price

Our Plan:

• 1. Compute max price in a subquery for a given company

S mot(prize) W cidel

Witnesses: with joins (2/6)

Product2 (pname, price, cid) Company2 (<u>cid</u>, cname, city)

Q: For each company, find the most expensive product + its price

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Our Plan:

• 1. Compute max price in a subquery for a given company

1. SELECT max(P1.price) FROM Product2 P1 WHERE P1.cid = 1 Witnesses: with joins (3/6)





Q: For each company, find the most expensive product + its price

Our Plan:

- 1. Compute max price in a subquery for a given company
- 2. Compute each product and its price, per company

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> 1. SELECT max(P1.price) FROM Product2 P1 WHERE P1.cid = 1

Witnesses: with joins (3/6)





Q: For each company, find the most expensive product + its price

Our Plan:

- 1. Compute max price in a subquery for a given company
- 2. Compute each product and its price, per company

2. SELECT C2.cname, P2.pname, P2.price FROM Company2 C2, Product2 P2 WHERE C2.cid = P2.cid

Witnesses: with joins (3/6)





Q: <u>For each company</u>, find the most expensive product + its price

Our Plan:

- 1. Compute max price in a subquery for a given company
- 2. Compute each product and its price, per company
- 3. Compare the price with the max price

2. SELECT C2.cname, P2.pname, P2.price FROM Company2 C2, Product2 P2 WHERE C2.cid = P2.cid

Witnesses: with joins (4/6)





Q: <u>For each company</u>, find the most expensive product + its price

Our Plan:

- 1. Compute max price in a subquery for a given company
- 2. Compute each product and its price, per company
- 3. Compare the price with the max price

```
SELECT C2.cname, P2.pname, P2.price

FROM Company2 C2, Product2 P2

WHERE C2.cid = P2.cid

and P2.price =

(SELECT max(P1.price)

FROM Product2 P1

WHERE P1.cid = C2.cid)
```

How many aliases do we actually need?

Witnesses: with joins (5/6)





Q: <u>For each company</u>, find the most expensive product + its price

Our Plan:

- 1. Compute max price in a subquery for a given company
- 2. Compute each product and its price, <u>per company</u> and compare the price with the max price

```
SELECT cname, pname, price
FROM Company2, Product2
WHERE Company2.cid = Product2.cid
and price =
(SELECT max(price)
FROM Product2
WHERE cid = Company2.cid)
```

We need no single alias here.

Next: can we eliminate the max operator in the inner query? Witnesses: with joins (6/6)





Q: <u>For each company</u>, find the most expensive product + its price

Our Plan:

- 1. Compute all prices in a subquery, for a given company
- 2. Compute each product and its price, <u>per company</u> and compare the price with the all prices

```
SELECT cname, pname, price
FROM Company2, Product2
WHERE Company2.cid = Product2.cid
and price >= ALL
(SELECT price
FROM Product2
WHERE cid = Company2.cid)
```

But: "ALL" does not work in SQLite ⊗ Product2 (pname, price, cid) Company2 (<u>cid</u>, cname, city)

Q: For each company, find the most expensive product + its price

Another Plan:

- 1. Create a table that lists the max price for each company id
- 2. Join this table with the remaining tables

F

 SELECT cid, max(price) as MP
 FROM Product2
 GROUP BY cid
 X, P, C

Finding the maximum price for each company was easy. But we want the "witnesses", i.e. the products with max price.



Q: For each company, find the most expensive product + its price

Another Plan:

- 1. Create a table that lists the max price for each company id
- 2. Join this table with the remaining tables

2. SELECT C2.cname, P2.pname, X.MP FROM Company2 C2, Product2 P2, (SELECT cid, max(price) as MP FROM Product2 GROUP BY cid) as X WHERE C2.cid = P2.cid and C2.cid = X.cid and P2.price = X.MP

Let's write the same query with a "WITH" clause





Q: For each company, find the most expensive product + its price

Another Plan with WITH:

- 1. Create a table that lists the max price for each company id
- 2. Join this table with the remaining tables

WITH X(cid, MP) as (SELECT cid, max(price)) FROM Product2 GROUP BY cid)
SELECT C2.cname, P2.pname, X.MP Company2 C2, Product2 P2, X
WHERE C2.cid = P2.cid and C2.cid = X.cid and P2.price = X.MP

#### Witnesses: with aggregates per group (1/8)



First: How to get the product that is sold with maximum price?

#### **Purchase**

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10

Product	mp
Banana	4

SELECT product, max(price) as mp FROM WHERE GROUP BY HAVING

???

### Witnesses: with aggregates per group (2/8)



*First: How to get the product that is sold with maximum price?* **Purchase** *1) Find the maximum price* 

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10





#### Witnesses: with aggregates per group (3/8)



First: How to get the product that is sold with maximum price? **Purchase** 2) Now you need to find product with price = maximum price

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10

Product	mp
Banana	4

SELECT	P2.product, P2.price as mp
FROM	Purchase P2
WHERE	P2.price = (
	SELECT max(price)
	FROM Purchase

#### Witnesses: with aggregates per group (4/8)



First: How to get the product that is sold with maximum price? **Purchase** Another way to formulate this query

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10

Product	mp
Banana	4

SELECT	P2.product, P2.price as mp		
FROM	Purchase P2		
WHERE	P2.price >= ALL (		
	SELECT price		
	FROM Purchase		

#### Witnesses: with aggregates per group (5/8)



Second: How to get the product that is sold with max <u>sales (=quanities sold)</u>? **Purchase** 

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10

	Product	sales
$\checkmark$	Banana	70

SELECT FROM WHERE GROUP BY HAVING	???

#### Witnesses: with aggregates per group (6/8)



Second: How to get the product that is sold with max <u>sales (=quanities sold)</u>? **Purchase** 1: find the total sales (sum of quantity) for each product

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10

Product	sales
Bagel	40
Banana	70

SELECTproduct, sum(quantity) as salesFROMPurchaseGROUP BY product

#### Witnesses: with aggregates per group (7/8)



Second: How to get the product that is sold with max <u>sales</u>? **Purchase** 2: we can use the same query as nested query

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10



# Witnesses: with aggregates per group (8/8)



Second: How to get the product that is sold with max <u>sales</u>? **Purchase** 3: putting things together

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10

$\sim$	Product	sales
$\mathcal{V}$	Banana	70

SELECT	product, sum(quantity) as sa	ales
FROM	Purchase	
<b>GROUP BY</b>	product	
HAVING	<pre>sum(quantity) &gt;= ALL (</pre>	
	SELECT sum(quantity)	
	FROM Purchase	
	GROUP BY product	)

Next: Can you write the query without the "ALL" quanitfier?

#### Witnesses: with aggregates per group (8/8)



Second: How to get the product that is sold with max <u>sales</u>? **Purchase** Another way to formulate this query without "ALL"

Product	Price	Quantity
Bagel	3	20
Bagel	2	20
Banana	1	50
Banana	2	10
Banana	4	10



SELECT proc	luct, sum(quantity) as sales
FROM Purc	hase
<b>GROUP BY</b> prod	uct
HAVING sum	(quantity) =
(SEI	ECT max (Q)
FRC	M (SELECT sum(quantity) Q
	FROM Purchase
	GROUP BY product) X

Understanding nested queries

### More SQL Queries

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid</u>, <u>bid</u>, <u>day</u>) Boats (<u>bid</u>, bname, color)



sid	sname	rating	age
22	Dustin	7	45.0
29	Brutus	1	33.0
31	Lubber	8	55.5
32	Andy	8	25.5
<b>5</b> 8	Rusty	10	35.0
64	Horatio	7	35.0
71	Zorba	10	16.0
74	Horatio	9	35.0
85	Art	3	25.5
95	Bob	3	63.5

Figure 5.1	An	Instance	S3	of	Sailors
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sid	bid	day
22	101	10/10/98
22	102	10/10/98
22	103	10/8/98
22	104	10/7/98
31	102	11/10/98
31	103	11/6/98
31	104	11/12/98
64	101	9/5/98
64	102	9/8/98
74	103	9/8/98

bid	bname	color
101	Interlake	blue
102	Interlake	red
103	Clipper	green
104	Marine	red

#### Figure 5.3 An Instance B1 of Boats

Figure 5.2 An Instance R2 of Reserves

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid</u>, <u>bid</u>, <u>day</u>) Boats (<u>bid</u>, bname, color)



Q: Find the names of sailors who have reserved a red boat.

SELECT S.sname FROM Sailors S WHERE S.sid IN (SELECT R.sid FROM Reserves R WHERE R.bid IN (SELECT B.bid FROM Boats B WHERE B.color = 'red'))

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Q: Find the names of sailors who have reserved a boat that is not red.

SELECT S.sname FROM Sailors S WHERE S.sid IN (SELECT R.sid FROM Reserves R WHERE R.bid not IN (SELECT B.bid FROM Boats B WHERE B.color = 'red'))

They must have reserved at least one boat in another color

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid</u>, <u>bid</u>, <u>day</u>) Boats (<u>bid</u>, bname, color)



Q: Find the names of sailors who have not reserved a red boat.

SELECT S.sname FROM Sailors S WHERE S.sid not IN (SELECT R.sid FROM Reserves R WHERE R.bid IN (SELECT B.bid FROM Boats B WHERE B.color = 'red'))

They can have reserved <u>0 or more boats</u> in another color, but must not have reserved any red boat

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Find the names of sailors who have reserved only red boatsQ: Find the names of sailors who have not reserved a boat that is not red.

SELECT S.sname FROM Sailors S WHERE S.sid not IN (SELECT R.sid FROM Reserves R WHERE R.bid not IN (SELECT B.bid FROM Boats B WHERE B.color = 'red'))

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Find the names of sailors who have reserved all red boatsQ: Find the names of sailors so there is no red boat that is not reserved by him.



To understand semantics of nested queries, think of a nested loops evaluation: For each Sailors tuple, check the qualification by computing the subquery

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Q: Find the names of sailors who have reserved a red boat.



Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Q: Find the names of sailors who have reserved a boat that is not red.



Query from: Ramakrishnan, Gehrke: Database management systems, 2nd ed (2000)

Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Q: Find the names of sailors who have not reserved a red boat.



Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Find the names of sailors who have reserved only red boatsQ: Find the names of sailors who have not reserved a boat that is not red.



Sailors (<u>sid</u>, sname, rating, age) Reserves (<u>sid, bid, day</u>) Boats (<u>bid</u>, bname, color)



Find the names of sailors who have reserved all red boatsQ: Find the names of sailors so there is no red boat that is not reserved by him.



# http://queryviz.com





Multiset operations (Intersect, Except)

# Recall Multisets (Bags)

Multiset X

Tuple
(1, a)
(1, a)
(1, b)
(2, c)
(2, c)
(2, c)
(1, d)
(1, d)



Equivalent Representations of a <u>Multiset</u>  $\lambda(X)$ = "Count of tuple in X" (Items not listed have implicit count 0)

#### Multiset X

Tuple	$\lambda(X)$
(1, a)	2
(1, b)	1
(2, c)	3
(1, d)	2

Note: In a set all counts are {0,1}.

# Generalizing Set Operations to Multiset Operations

#### Multiset X

Tuple	$\lambda(X)$
(1, a)	2
(1, b)	0
(2 <i>,</i> c)	3
(1, d)	0

#### Multiset Y

Tuple	$\lambda(Y)$
(1, a)	5
(1, b)	1
(2 <i>,</i> c)	2
(1, d)	2

#### Multiset Z

Tuple	$\lambda(Z)$
(1, a)	2
(1, b)	0
(2 <i>,</i> c)	2
(1, d)	0

$$\lambda(Z) = min(\lambda(X), \lambda(Y))$$

For sets, this is intersection

# Generalizing Set Operations to Multiset Operations

#### Multiset X

Tuple	$\lambda(X)$
(1, a)	2
(1, b)	0
(2, c)	3
(1, d)	0

#### Multiset Y

Tuple	$\lambda(Y)$
(1, a)	5
(1, b)	1
(2, c)	2
(1, d)	2

Multiset Z

Tuple	$\lambda(Z)$
(1, a)	7
(1, b)	1
(2, c)	5
(1, d)	2

$$\lambda(Z) = \lambda(X) + \lambda(Y)$$

For sets, this is **union** 

# Multiset Operations in SQL

# Explicit Set Operators: INTERSECT

$$\{r.A \mid r.A = s.A\} \cap \{r.A \mid r.A = t.A\}$$

UNION

 ${r.A | r.A = s.A} \cup {r.A | r.A = t.A}$ 



Why aren't there duplicates?

By default: SQL uses set semantics for INTERSECT and UNION!

What if we want duplicates?

### UNION ALL

 ${r.A | r.A = s.A} \cup {r.A | r.A = t.A}$ 



ALL indicates Multiset operations EXCEPT

 $\{r.A \mid r.A = s.A\} \setminus \{r.A \mid r.A = t.A\}$ 



What is the multiset version?

### INTERSECT and EXCEPT\*





If R, S have no duplicates, then can write without sub-queries (HOW?)



\*Not in all DBMSs. (SQLlite does not like the parentheses, Oracle uses "MINUS" instead of "EXCEPT")