CS3000: Algorithms & Data — Summer '23 — Laney Strange

Mergesort Algorithm

We went over the divide-and-conquer Mergesort algorithm in class; this handout is so you can see the function typeset in the CLRS pseudocode style. The algorithm *divides* by making the input array smaller and smaller and smaller until it's trivially easy to sort (sorting the trivially easy is the *conquer* step). Then it merges together the smaller sorted subarrays – that's the *combine* step.

```
MERGESORT(A, p, q, r)
1  if p < r
2        q = ⌊(p + r)/2⌋
3        MERGESORT(A, p, q)
4        MERGESORT(A, q + 1, r)
5        MERGE(A, p, q, r)
```

```
MERGE(A, p, q, r)
 1  n_L = q − p + 1
 2  n_R = r − q
 3  let L[1 : n_L + 1] and R[1 : n_R + 1] be new arrays
 4  for i = 1 to n_L
 5        L[i] = A[p + i − 1]
 6  for j = 1 to n_R
 7        R[j] = A[q + j]
 8  L[n_L + 1] = ∞
 9  R[n_R + 1] = ∞
10  i = 1
11  j = 1
12  k = p
13  while k ≤ r
14        if L[i] ≤ R[j]
15              A[k] = L[i]
16              i = i + 1
17        else
18              A[k] = R[j]
19              j = j + 1
20        k = k + 1
```

Here's how we typeset the above functions in CLRS style:

```
\begin{codebox}
\Procname{$\proc{Mergesort}(A, p, r)$}
\li \If $p \ge r$
\Then
\li \Return
```

```
\li $q = \lfloor(p+r)/2\rfloor$
\li $\proc{Mergesort}(A, p, q)$
\li $\proc{Mergesort}(A, q+1, r)$
\li $\proc{Merge}(A, p, q, r)$
\end{codebox}

\medskip


\begin{codebox}
\Procname{$\proc{Merge}(A, p, q, r)$}
\li $n_L = q - p + 1$
\li $n_R = r-q$
\li let $L[1:n_L + 1]$ and  $R[1:n_R + 1]$ be new arrays
\li \For $i \gets 1 \To n_L$
\Do
\li $L[i] = A[p+i-1]$
\End
\li \For $j \gets 1 \To n_R$
\Do
\li $R[j] = A[q+j]$
\End
\li $L[n_L + 1] = \infty$
\li $R[n_R + 1] = \infty$
\li $i = 1$
\li $j = 1$
\li $k = p$
\li \While $k \le r$
\Do
\li \If $L[i] \le R[j]$
\Then
\li $A[k] = L[i]$
\li $i = i + 1$
\li \Else
\li $A[k] = R[j]$
\li $j = j + 1$
\End
\li $k = k + 1$
\End
\end{codebox}
```