CS3000: Algorithms & Data — Summer 2023 — Laney Strange

Recitation 2 Due May 16th @ 9pm via Gradescope

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- One recitation problem each week is graded; the rest are there for practice. That problem will be graded on completeness full credit for making an honest effort. It is also closely linked to the quicksort problem on the upcoming Short Homework, so it's good to review it now!
- Recitations can be written by hand; submit a picture/screenshot on Gradescope.
- Put your name on the first page. If you are using the LATEX template we provided, then you can make sure it appears by filling in the yourname command.
- This recitation is due May 16th @ 9pm via Gradescope. If you miss the in-person recitation, or need to submit your solution later than the end of your section, please fill out this form: https://forms.gle/CLrhrkVauXYzC7U57
- Collaboration is strongly encouraged during recitation! Please list all your collaborators in your solution for each problem by filling in the yourcollaborators command.

Problem 1. Quicksort (graded)

(a) You can find the pseudocode for the Quicksort algorithm we did in class at https://course.ccs.neu.edu/cs3000/resources/Quicksort_Pseudocode.pdf. What is the runtime of the Partition step of Quicksort?

Solution:

(b) Give an example of an array that would result in the worst-case run-time of Quicksort.

Solution:

(c) Give a recurrence for the worst-case run-time of Quicksort.

Solution:

(d) Give an example of an array that would result in the best-case run-time of Quicksort.

Solution:

(e) Give a recurrence for the best-case run-time of Quicksort

Solution:

Problem 2. Karatsuba (for practice; not graded)

You can find the pseudocode we did in class for the Karatsuba algorithm at https://course.ccs.neu.edu/cs3000/resources/Karatsuba_Pseudocode.pdf

Let's walk through how the algorithm works when multiplying 7832 · 3671.

(a) The first call to Karatsuba is KARATSUBA(7832, 3671, 4). What are *m*, *a*, *b*, *c*, *d*?

Solution:

(b) The second call to Karatsuba is ac = KARATSUBA(78, 36, 2). What are *m*, *a*, *b*, *c*, *d*?

Solution:

(c) In the middle of that second call, we make three more calls to Karatsuba, but they will all end up in the base case. What are *ac*, *bd*, *bacd* after returning from the base case calls?

Solution:

(d) Using those results, what does the second call return?

Solution:

(e) After all the recursive calls have completed, we're back in the first call to Karatsuba, like in Part A above. After all the recursive calls have completed, we now have *ac* = 2808, *bd* = 2272, *bacd* = 1610. What is the final answer?

Solution:

Problem 3. *Mergesort (for practice; not graded)*

Suppose you start with the following array: < 8, 5, -9, 14, 0, -1, -7, 3 >

(a) How many subarrays would mergesort split this array into in its first step? What would they be?

Solution:

(b) How many subarrays would mergesort split this array into in its second step? What would they be?

Solution:

(c) How many subarrays would mergesort split this array into in its third step? What would they be?

Solution:

(d) How many subarrays would exist after the first merge step? What would they be?

Solution:

(e) How many subarrays would exist after the second merge step? What would they be?

Solution:

(f) How many subarrays would exist after the third merge step? What would they be?

Solution:

(g) We know that the bound on Mergesort's worst-case run-time is $\Theta(n \lg n)$. What is the bound on its best-case run time?

Solution:

(h) Suppose you start with an already-sorted array. Describe a small change you can make to the Mergesort algorithm that could take advantage of this for a better real-world runtime.

Solution: