

CS3000

5/11 - Thurs 'i

Admin

- no rec today (first thus func is 5/18)
- Shift HW1 due 9pm
- Long ~~HW~~ HW1 out, due Tues. 9pm

Agenda

1. Divide and conquer on a number
2. Karatsuba algorithm
3. Algorithm Analysis

Math Recap

$$\text{(ex)} \quad \sum_{i=1}^n i = \frac{(n)(n+1)}{2} = \Theta(n^2)$$

$$\text{(ex)} \quad \sum_{i=0}^n x^i = \frac{x^{n+1} - 1}{x - 1} = \Theta(x^n)$$

$$\text{(ex)} \quad 2^{\lg c} = c^{\lg 2} \quad c=8 \quad 2=4$$

$$4^{\lg 8} = 4^3 = 8^{\lg 4} = 8^2 = 64$$

$$\text{(ex)} \quad 2^{\lg n} = n \quad 2^{\lg n - 1} = \frac{n}{2}$$

1. Divide + Conquer on 2 number

Linear Search \rightarrow Bin Search

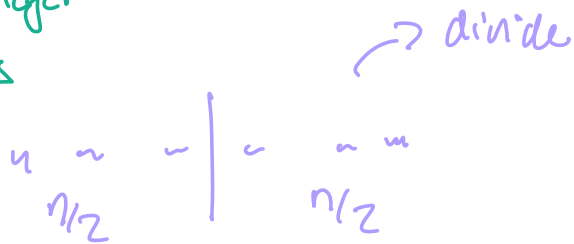
$\Theta(n)$

$\Theta(\lg n)$

W.C.

\hookrightarrow divide + conquer on BST

Today: input is an integer with n digits



Using a number with n digits in Multiplication

(ex) Addition first

$n=4$

$w = 1\ 2\ 3\ 4$

$y = 1\ 1\ 2\ 2$

$w+y$

$$\begin{array}{r}
 1\ 2\ 3\ 4 \\
 +\ 1\ 1\ 2\ 2 \\
 \hline
 2\ 3\ 5\ 6
 \end{array}$$

$\Theta(n)$ algorithm

(ex) Multiplication

$w = 1\ 2\ 3\ 4$

$y = 1\ 1\ 2\ 2$

$w \times y$

$$\begin{array}{r}
 1\ 2\ 3\ 4 \\
 \times\ 1\ 1\ 2\ 2 \\
 \hline
 2\ 4\ 6\ 8 \\
 2\ 4\ 6\ 8\ 0
 \end{array}$$

$\Theta(n^2)$

$$\begin{array}{r} \\ \\ \hline 1 \\ 1 \\ 2 \\ 3 \\ 4 \\ 0 \\ 0 \\ 0 \end{array}$$

Can we do better?

↳ maybe with 2 ^{known} algorithmic techniques!

Insights

- adding on zeros is easy
 - ↳ multiply by a power of 10
 - do not real multiplication
- express a decimal number with powers of 10
 - ↳ $1234 = 1 \cdot 10^3 + 2 \cdot 10^2 + 3 \cdot 10^1 + 4 \cdot 10^0$
 - ↳ $1234 = 12 \cdot 10^2 + 34 \cdot 10^0$
 - $= 1200 + 34$
 - $= 1234$

$$W = 1234 = 12 \cdot 10^2 + 34 \cdot 10^0$$
$$Y = 1122 = 11 \cdot 10^2 + 22 \cdot 10^0$$



→ want to keep multi to $n/2$ digits

$$W \times Y = 1234 \times 1122$$

$$= (12 \cdot 10^2 + 34)(11 \cdot 10^2 + 22)$$

$$= 10^4 \cdot 12 \cdot 11 + 34 \cdot 11 \cdot 10^2 + 22 \cdot 12 \cdot 10^2 + 34 \cdot 22$$

$$= 10^4 \cdot 12 \cdot 11 + 10^2(34 \cdot 11 + 22 \cdot 12) + 34 \cdot 22$$

$\xrightarrow{\text{red}} \hookrightarrow n/2 \text{ mult.}$ $\xrightarrow{\text{red}} \hookrightarrow n/2 \text{ mult.}$

$$= 1384548$$

∩

? Actually faster?

2. Karatsuba Algorithm

$$n=4$$

$$W = \frac{12}{a} \frac{34}{b}$$

$$Y = \frac{11}{c} \frac{22}{d}$$

$$W = 10^{n/2} \cdot a + b$$

$$Y = 10^{n/2} \cdot c + d$$

$$W \times Y = (10^{n/2} \cdot a + b)(10^{n/2} \cdot c + d)$$

$$= 10^n \cdot a \cdot c + 10^{n/2} \cdot b \cdot c + 10^{n/2} \cdot a \cdot d + b \cdot d$$

$$= 10^n \cdot a \cdot c + 10^{n/2}(b \cdot c + a \cdot d) + b \cdot d$$

∩

The Algo

KAR(w, y, n)

if $n = 1$

→ conquer

return $w \cdot y$

$m = \lceil n/2 \rceil$

// Split w into z, b and y into c, d

$e = \text{KAR}(z, c, m)$

$z \cdot c$

→ divide

$f = \text{KAR}(b, d, m)$

$b \cdot d$

$g = \text{KAR}(b-z, c-d, m)$

$(b-z)(c-d)$

return $10^{2m}e + 10^m(e + f + g) + f$

→ combine

ex) $w = 48$ $n = 2$
 $y = 34$

$48 \cdot 34 = ? \quad 1632 \quad !!$

↳ $m = 1$

Split $w = \begin{matrix} 4 & 8 \\ \hline z & b \end{matrix}$ $y = \begin{matrix} 3 & 4 \\ \hline c & d \end{matrix}$

$e = \text{KAR}(4, 3, 1) \rightarrow 12$

$f = \text{KAR}(8, 4, 1) \rightarrow 32$

$g = \text{KAR}(4, 1, 1) \rightarrow -4$

return $10^2 \cdot 12 + 10(12 + 32 - 4) + 32$

↳ $1200 + 400 + 32$

↳ 1632

$10:47$

3. Algo Analysis

- Run Time
- Correctness

$$T(n) = \dots ??$$

```
KAR(w, y, n)
if n == 1
    return w * y
```

Base case $T(1) = 2$

$$m = \lceil n/2 \rceil$$

C

// Split w into z, b and y into c, d

$$e = KAR(z, c, m)$$

$$f = KAR(b, d, m)$$

$$g = KAR(b-z, c-d, m)$$

$3 \cdot T(n/2)$

$$\text{return } 10^{2m} e + 10^m (e + f + g) + f \quad C \cdot n$$

$$T(n) = 3 \cdot T(n/2) + cn + C$$

$$T(1) = 2$$

} run time.
need to solve it!

Iteration method

1st $T(n) = 3 \cdot T(n/2) + cn + C$

Plug in: $T(n/2) = 3 \cdot T(n/4) + C \cdot \frac{n}{2} + C$

2nd $T(n) = 3 [3 \cdot T(n/4) + \frac{1}{2}cn + C] + cn + C$

$$= 9 T(n/4) + \frac{3}{2}cn + cn + 3c + c$$

Plug in: $T(n/4) = 3 \cdot T(n/8) + c \cdot \frac{n}{4} + c$

(3rd) $T(n) = 9 \left[3 \cdot T(n/8) + \frac{1}{4}cn + c \right] + \frac{3}{2}cn + cn + 3c + c$

$$= 27 T(n/8) + \underbrace{\frac{9}{4}cn + \frac{3}{2}cn + cn}_{\substack{3^2 \\ 2^2} cn \quad \substack{3^1 \\ 2^1} cn \quad \substack{3^0 \\ 2^0} cn} + \underbrace{9c + 3c + c}_{\substack{3^2 c \quad 3^1 c \quad 3^0 c}}$$

(kth)

$$\sum_{i=0}^{k-1} \left(\frac{3}{2}\right)^i \cdot cn$$

$$cn \sum_{i=0}^{k-1} \left(\frac{3}{2}\right)^i$$

$$cn \cdot \Theta\left(\frac{3}{2}\right)^{k-1}$$

$$\sum_{i=0}^{k-1} 3^i \cdot c$$

$$c \cdot \sum_{i=0}^{k-1} 3^i$$

$$c \cdot \Theta\left(3\right)^{k-1}$$

→ Bounds on sums

(cn)

$$T(n) = 3^k \cdot T(n/2^k) + cn \left(\frac{3}{2}\right)^{k-1} + c \cdot (3)^{k-1}$$

(choose a value for k to get to base case $T(1) = 2$)

$$k = \lg n$$

$$T(n) = 3^{\lg n} \cdot T(n/n) + cn \left(\frac{3^{\lg n - 1}}{2^{\lg n - 1}}\right) + c \cdot 3^{\lg n - 1}$$

$$= 3^{\lg n} \cdot 2 + \frac{cn \cdot 2}{2} \cdot (3^{\lg n - 1}) + c \cdot 3^{\lg n - 1}$$

$$= 3^{lg n} \cdot 2 + 2c \cdot (3^{lg n - 1}) + c \cdot 3^{lg n - 1}$$

$$= \Theta(3^{lg n})$$

$$= \Theta(n^{lg 3}) = \Theta(n^{1.59})$$

→ Beats $\Theta(n^2)$!

Proof of correctness → by induction on n
(n is #digits)

$$w \cdot y = 10^n \cdot a \cdot c + 10^{n/2} (b \cdot c + a \cdot d) + b \cdot d$$

Base case: $n=1$

↳ KAR returns $w \cdot y$. done!

(IH) KAR works correctly when w, y have
1 digit, 2 digits, ..., $k-1$ digits

↳ strong ind.

Want to show it works with (k) digits

In algo...
 $m = k/2$

$$e = \text{KAR}(a, c, m)$$

$$f = \text{KAR}(d, b, m)$$

$$g = \text{KAR}(b-a, c-d, m)$$

$$\text{return } 10^{2m} \cdot e + 10^m (e + f + g) + f$$

these work!

By ind. Hyp!

$$10^k \cdot (a \cdot c) + 10^{k/2} (zc + bd + (b-a)(c-d)) + bd$$

$$= 10^k zc + 10^{k/2} (zc + bd + bc - zc - bd + zd) + bd$$

$$\boxed{= 10^k zc + 10^{k/2} (bc + zd) + bd}$$

✓ done!
!!
☺

• Algo's runtime is $\Theta(n^{1.59})$

↳ better: $\Theta(n^2)$

• Algo's correctness with induction

