

CS3000  
6/20-Tues.

### Admin

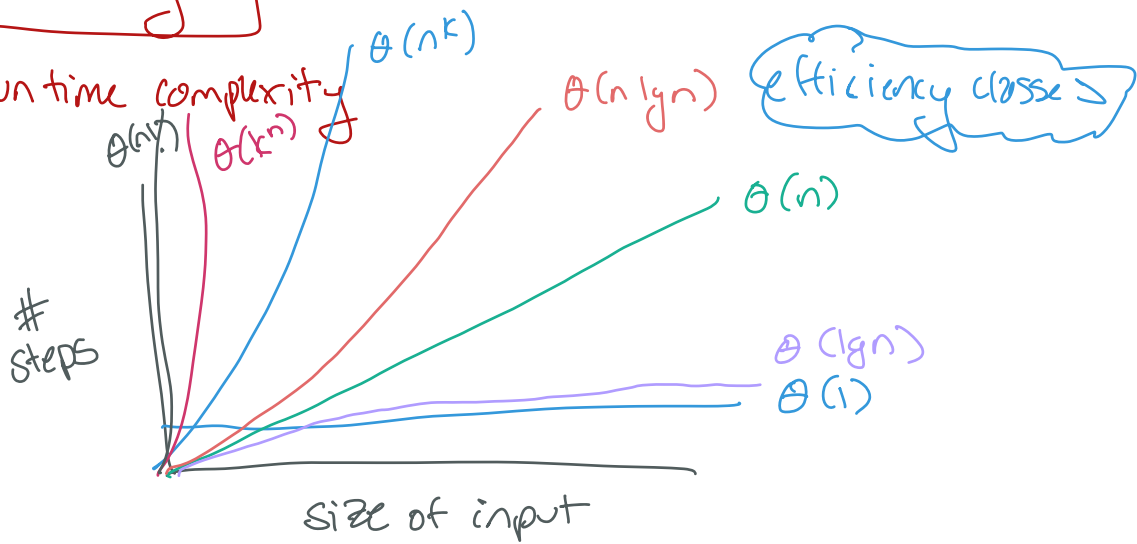
- exam back today / solutions
- make-up exam question 6/22
- Recitation 5 today → last recitation !!  
6/20 9pm → submit Rec 3-5 } no extensions
- Second trace hw due 6/20 9pm
- Short HW4 at now due 6/22 9pm
- Form for tomorrow's topics
- Please do TRACE !!

### Agenda

1. Interact2bility
  2. P vs. NP
-

# 1. Intractability

↳ runtime complexity



$\theta(1)$

$\theta(\lg n)$

$\theta(n)$  ( $\theta(V+E)$ )

$\theta(n \lg n)$  (comparison based sorting  $\Omega(n \lg n)$ )

$\theta(n^k)$  (polynomial)  $\theta(V \cdot E)$

tractable !!

$\theta(k^n)$  exponential

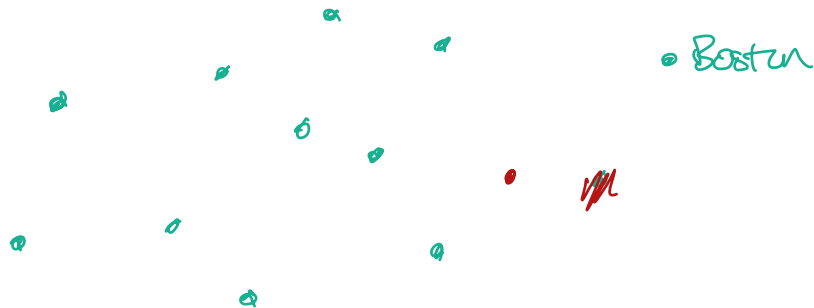
$\theta(n!)$  factorial

intractable !!

↳ may be a solution, but can't really run on a computer

# Traveling Salesman Problem (TSP)

person



Salesman starts in Boston  
wants to visit n cities  $\rightarrow$  optimal way

like SSSP/APSP but visit all n before home

• only known solution is to visit all  $n!$  <sup>permutations</sup> ~~cities~~  
pick the best one

Run-time:  $\Theta(n!)$

What does  $n!$ , run-time mean?

• ex: computer does  $2^{16} = 65,536$  operations per second

$n$	$n!$	Wall-clock time
4	24	$24 / 65536 = .00036$
8	40320	.615 seconds
16	20,922,789,888,000	3695 days
32	?	$1.271 \times 10^{20}$ millennia

# TSP

- no known  $\Theta(n^k)$  or better solution
- But, don't know if one exists or not

## NP complete

millenial problem  
\$1m prize to solve either way

## 2. P vs NP

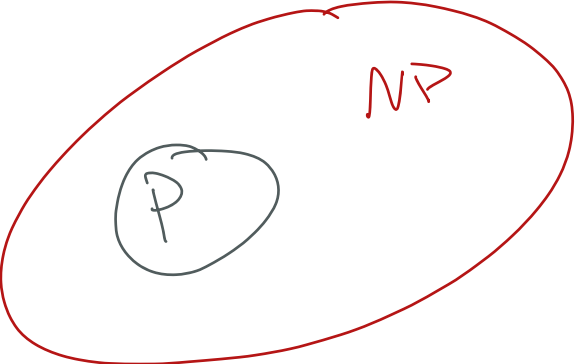
P - solvable in  $\Theta(n^k)$  time  
by a deterministic Turing machine



NP - solvable in  $\Theta(n^k)$  time  
by a non-deterministic Turing machine



(Brute force solution)



- Class P
- poly time soln
  - Brute force soln

- Class NP
- Brute force solution
  - poly time soln may exist
  - solution can be verified

in  $\Theta(n^4)$  time

(not all problems in CS belong to NP)

NP/P  $\rightsquigarrow$  decision problems

- possible answers: yes/no
- optimization  $\rightarrow$  decision

↳ ex. SSSP: is there a path from  $s$  to  $v$   
of length  $\leq l$ ?

- P vs. NP has big implications for optimization problems

P vs. NP problems

• SSSP  $\rightsquigarrow$   $\in P$

• SS LP (general graph)  $\rightsquigarrow$   $\in NP$   
 $\in P??$  no known  
poly time soln

• Eulerian path (edges)  $\rightsquigarrow$   $\in P$

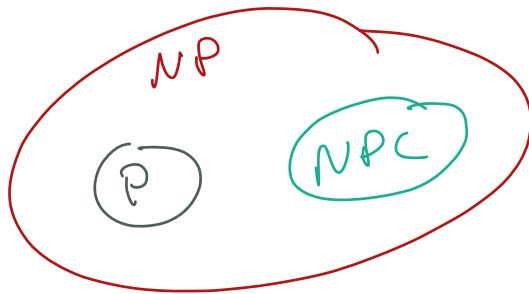
• Hamiltonian path (vertices)  $\rightsquigarrow$   $\in NP$   
 $\in P??$  no known  
poly time soln

NP complete:

• in NP

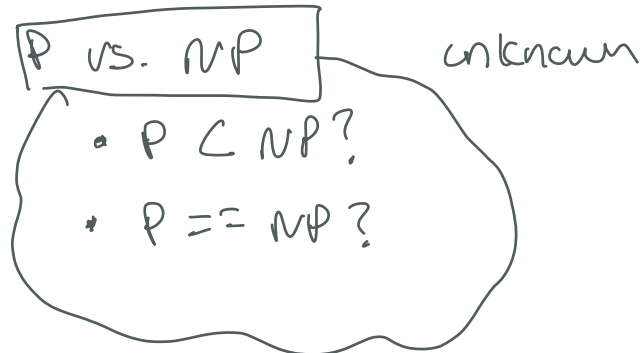
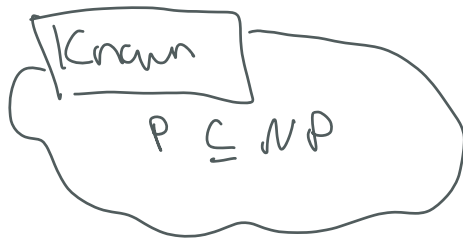
• not known to be in P

} TSP, SS LP, Hamilton



\$ question:

- what is relationship between P, NP?
- All NPC are related to each other



NP-hard

- as hard as any NP problem
- all NPC are NP-hard
- some NP-hard problems not in NP

