

CS3000

6/6 - Tues.

Admin

- Long HW4 due today 9pm
- Short HW3 out now, due Thurs 9pm
- Rec 4 today!
- Fun optional recitation Thurs

Agenda

1. Minimum Spanning Trees
2. Kruskal
3. Prim

Recap

DFS \rightarrow topological?

track the order in which we hit dead ends
(v.f)

DFS \rightarrow determine if there's a cycle?

if there's a back edge, there's a cycle

DFS run-time

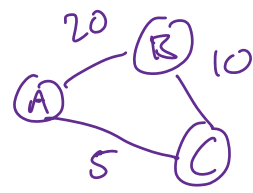
$\Theta(V+E)$ (pretty good!)

1. Minimum Spanning Trees

- apply techniques we know to graph problems

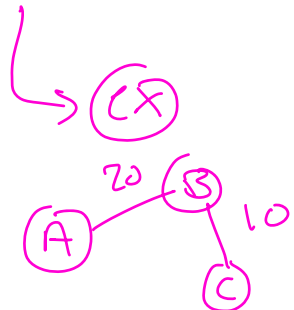
Today:

- undirected } graph G
- weighted }



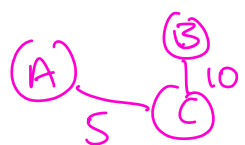
Want:

- Spanning tree
↳ tree subset of edges
include all vertices



- minimum spanning tree
↳ least total weight

tree ✓
all vertices ✓
total weight = 30



tree ✓
all vertices ✓
total weight = 15

MST

MST is an optimization problem

DP
Greedy ←

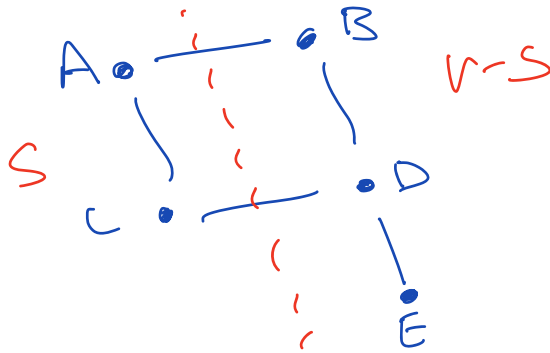
Implementation Detail

- weight function, parameter to all functions

Func (G, w) → $w(u, v)$ weight of edge (u, v)

Definitions

- MST is a set of edges
- Set $A \subseteq \text{MST}$ (while we're building)
- safe edge can be added to A , and A is still $\subseteq \text{MST}$
- A cut $(S, V-S)$ is a partition of vertices V



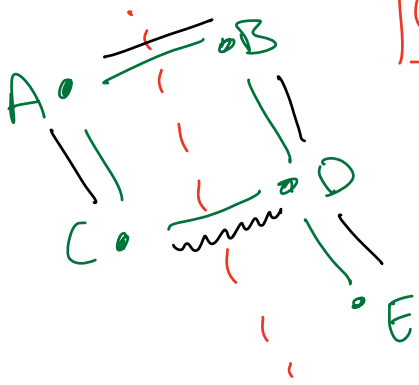
- An edge (u, v) with $u \in S$ and $v \in V-S$ "crosses the cut" above: (A, B) (C, D)
- A light edge is edge with min weight crossing the cut

Both algos...

- cut graph (keep some vertices in S)
- add light edge to MST (A)

↳ greedy choice!





Greedy works for MST

(C,D) is light edge

Suppose we choose (A,B) instead

↳ we can swap in ~~(A,B)~~
(C,D) and get a better tree

[adding in (C,D) will create a cycle in mst, we can remove some edge to get back to a tree.]

2. Kruskal + Prim

- graph G
- undirected
- weighted
- Adj list

Sets

- MAKE-SET(v) → create one-element set with v
- FIND-SET(v) → returns rep from v 's set
- $A \cup \{v\}$ → add element to set A
- UNION(u, v) → draw edge between u and v (puts u, v in same set)



Heaps

- $\text{INSERT}(H, v)$ \rightarrow insert v into H and heapify
- $\text{EXTRACT-MIN}(H)$ \rightarrow removes + return smallest element, and heapify
- $\text{DECREASEKEY}(H, v, w)$ \rightarrow replace v 's key with w , and heapify

Kruskal: make set of edges that form MST

Prim: track v .key and v . π on every vertex

What we want to know...

- Kruskal: at each step, what's in A , what are sets S
- Prim: at each step, what are v .key and v . π
- What is total weight of MST?
- Be ready to answer Qs from other team

W: 46

KRUSKAL(G, w)

```

1  A = {}
2  for each vertex v in G.V
3      MAKE-SET(v)
4  create a single list of the edges in G.E
5  sort the list of edges into monotonically increasing order by weight w
6  for each edge (u, v) taken from the list in sorted order
7      if FIND-SET(u) ≠ FIND-SET(v)
8          A = A ∪ {(u, v)}
9          UNION(u, v)
10 return A

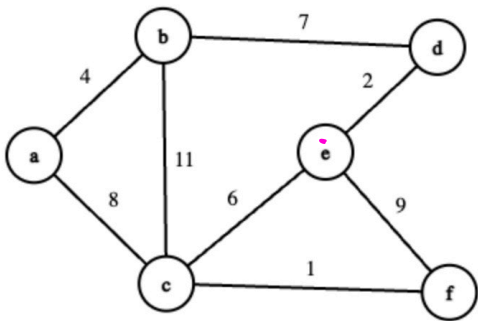
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PRIM(G, w, r)

```

1  for each vertex u in G.V
2      u.key = ∞
3      u.π = NIL
4  r.key = 0
5  H = {}
6  for each vertex u in G.V
7      INSERT(H, u)
8  while H ≠ {}
9      u = EXTRACT-MIN(H)
10     for each vertex v in G.adj[u]
11         if v in H and w(u, v) < v.key
12             v.π = u
13             v.key = w(u, v)
14     DECREASE-KEY(H, v, w(u, v))

```



Kruskal

A $\{\}$

Sets a, b, c, d, e, f

Step #1

(c, f) weight 1

$A = \{(c, f)\}$

Sets: a, b, d, e, c, f

Step #2

(e, d) w 2

$A = \{(c, f), (e, d)\}$

Sets: a, b, d, e, c, f

Step #3

(a, b) w 4

$A = \{(c, f), (e, d), (a, b)\}$

Sets: a, b, d, e, c, f

Step #4

(c, e) w 6

$A = \{(c, f), (e, d), (a, b), (c, e)\}$

Sets: a, b, d, e, c, f

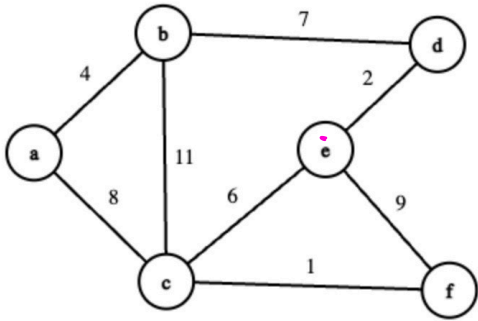
Step #5

(b, d) w 7

$A = \{(c, f), (e, d), (a, b), (c, e), (b, d)\}$

Sets: a, b, d, e, c, f

1 2 4 6 7 = 20



Prim

	A	B	C	D	E	F
key	0	∞	∞	∞	∞	∞
π	nil	nil	nil	nil	nil	nil

Step #1

H = A (smallest)

	A	B	C	D	E	F
key	0	4	8	∞	∞	∞
π	nil	A	A	nil	nil	nil

Step #2

extract B

H = BCDEF

	A	B	C	D	E	F
key	0	4	8	7	∞	∞
π	nil	A	A	B	nil	nil

Step #3

extract D

H = CDEF

	A	B	C	D	E	F
key	0	4	8	7	2	∞
π	nil	A	A	B	D	nil

Step#4

Extract E

H = CEF

	A	B	C	D	E	F
key	0	4	6	7	2	9
π	nil	A	E	B	D	E

Step#5

Extract C

H = CF

	A	B	C	D	E	F
key	0	4	6	7	2	9
π	nil	A	E	B	D	C

Step#6

Extract F

H = F

	A	B	C	D	E	F
key	0	4	6	7	2	9
π	nil	A	E	B	D	C

Final

	A	B	C	D	E	F
key	0	4	6	7	2	1
π	nil	A	E	B	D	C

total weight = $4 + 6 + 7 + 2 + 1 = 20$ omg !!

Run Time

PRIM(G, w, r)

```

1 for each vertex  $u \in G.V$ 
2    $u.key = \infty$ 
3    $u.\pi = NIL$ 
4  $r.key = 0$ 
5  $H = \{r\}$ 
6 for each vertex  $u \in G.V$ 
7   INSERT( $H, u$ )
8 while  $H \neq \{r\}$ 
9    $u = EXTRACT-MIN(H)$ 
10  for each vertex  $v \in G.adj[u]$ 
11    if  $v \in H$  and  $w(u, v) < v.key$ 
12       $v.\pi = u$ 
13       $v.key = w(u, v)$ 
14      DECREASE-KEY( $H, v, w(u, v)$ )
  
```

$v \lg v$

$V \times E$

$\lg v$ each time

$\Theta((V+E) \lg V)$

loop + heap

KRUSKAL(G, w)

```

1  $A = \{ \}$ 
2 for each vertex  $v \in G.V$ 
3   MAKE-SET( $v$ )
4 create a single list of the edges in  $G.E$ 
5 sort the list of edges into monotonically increasing order by weight  $w$ 
6 for each edge  $(u, v)$  taken from the list in sorted order
7   if FIND-SET( $u$ )  $\neq$  FIND-SET( $v$ )
8      $A = A \cup \{(u, v)\}$ 
9     UNION( $u, v$ )
10 return  $A$ 
  
```

E

Sorting: $\Theta(E \lg E)$
Band by sort!

(Roughly the same)