

CS3000

5/30 - Tue.

Admin

- Exams back today!
- HW3 - Long due tom. 9pm
- HW4 - Long goes out thurs
- Rec 3 today (graded)
- Fun optional recitation thurs.

Agenda

1. Heaps
2. Heapsort Overview
3. Heapsort Implementation

1. Heaps

Heapsort ~ sort of like selection sort

Find smallest unsorted element
 $\Theta(n^2)$

Can we do better?
//

} Apply a Strategy
Quicksort
mergesort } D+C

→ Another way to do better:

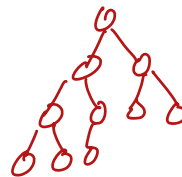
Change the data structure

- makes algo faster ←
- or, need/went to do the algorithm

Heap

A heap is a complete binary tree

- Binary tree: every node has ≤ 2 children
- Complete: levels get filled in left to right
- Height: $\Theta(\lg n)$



A min heap:

- every node is smaller than all its descendants

Heap Properties:

- min, completeness

Comparison

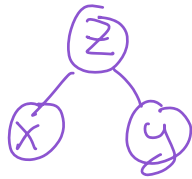
BT



Z.left = x
 Z.right = y

(tree imp.)

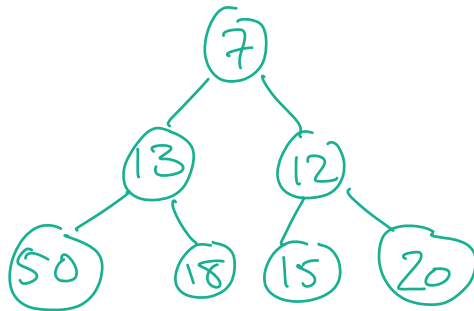
Heap



$\langle Z, X, Y \rangle$

(array imp.)

ex



min heap

Array: $\langle 7, 13, 12, 50, 18, 15, 20 \rangle$
 1 2 3 4 5 6 7

Q: For a node at i , where are its children?

left $2i$

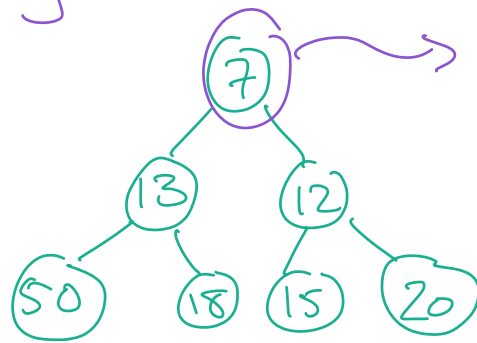
right $2i+1$

Array A to represent a Heap:

attribute A.heapsize (pos in A where heap ends)

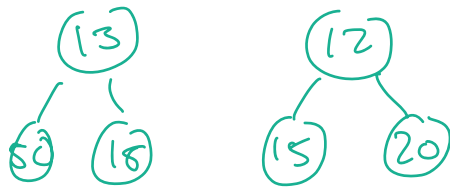
2. Heapsort Overview

- Given a min-heap, stored in array
- Take the min element (position 1), put in next space in output
- Re-heapify!
- Sort in descending order



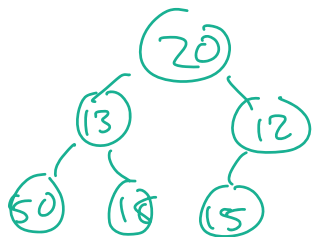
remove
root
(smallest)

Sorted
[7]



Heapify

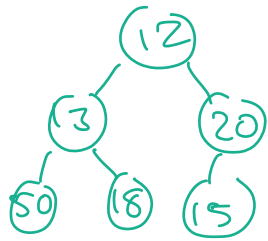
- complete BT
- min-ness
- last element is new root



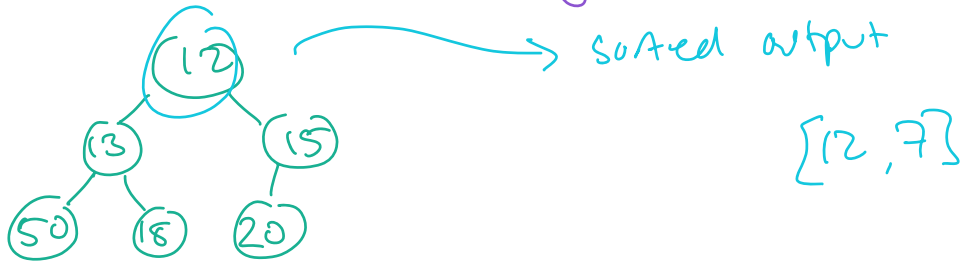
(maintains structure
need to fix min-ness)

mostly perfect min-heap
(root might be messed up)

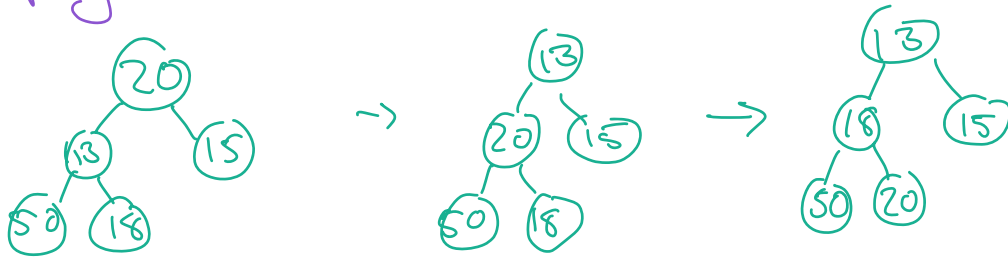
- swap 20 with 12 (swap root with smaller of 2 children)



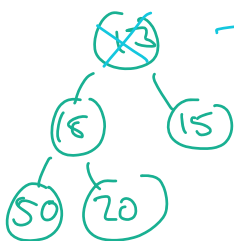
- swap 20 with 15 (only child)



- Heapify!



- remove smallest element



Sorted output
[13, 12, 7]

- Heapify



50

50

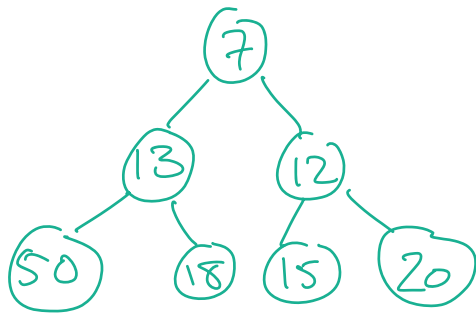
Run-time of ^{one}Heapify (worst case)

$\Theta(\log n)$

→ Finding the min element!

10:38

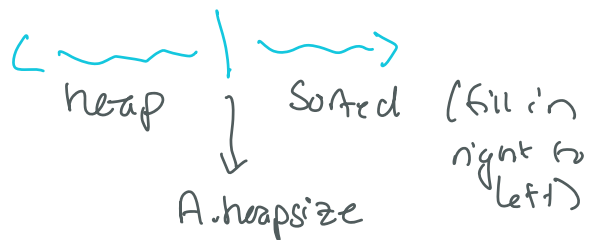
3. Heapsort Implementation



→ $\langle 7, 13, 12, 50, 18, 15, 20 \rangle$

• In-place algorithm

↳ at any point, some subarray is sorted, and some is still a heap



• Call to Heapify function

↳ today: black box
 takes in A, i
 ↳ current root

- assume the root is messed up, but the rest is ok
- bubbles down the root

Heapsort (A) $A \rightarrow$ starts as min heap

$A.heapsize = A.length$ \rightsquigarrow starts with whole thing unsorted

for $i = A.length$ down to 2

 swap $A[i], A[1]$ \rightarrow swap root with last ele

$A.heapsize = A.heapsize - 1$

 Heapify (A, 1) \rightarrow reheapify

> $\langle 7, 13, 12, 50, 18, 15, 20 \rangle$
 1 2 3 4 5 6 7

$A.heapsize = 7$
 $i = 7$
 swap $A[i], A[7]$
 $A.heapsize = 6$
 Heapify (A, 1)

↓
 swap

↓

$\langle 20, 13, 12, 50, 18, 15, 7 \rangle$
 1 2 3 4 5 6 7
 (heap) (sorted)

↓
 Heapify

< 12, 13, 15, 50, 18, 20, 7 >

↓
Swap

< 20, 13, 15, 50, 18, 12, 7 >
(heap) (sorted)

↓
Heapify

< 13, 18, 15, 50, 20, 12, 7 >
(Swap)

< 20, 18, 15, 50, 13, 12, 7 >
(heap) (sorted)

$i = 6$

Swap [1], [6]

A.heapsize = 5

Heapify (A, i)

$i = 5$

Swap [1], [5]

A.heapsize = 4

Heapify (A, i)

Runtime of Heapsort:

- loop: n times
- Heapify: $\lg n$ each call

↳ total: $\Theta(n \lg n)$

Did we beat selection sort?

Yes!

Did we take up extra space?

No!

