

CS3000

5/23 - Tue.

## Admin

- Long HW2 due 9pm
- today's recitation: unofficial
- Solns to practice problems are on piazza
- next HW goes out Thurs
- exam Thurs in class

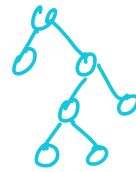
## Agenda

1. Greedy Strategy
2. Huffman codes
3. Greedy Huffman

## Recap

Binary tree every node has  $\leq 2$  children

Binary Search tree left == smaller  
right == bigger



What does a greedy algo do?  
makes locally best choice

Greedy choice for scheduling problem  
earliest finish time

## 1. Greedy Strategy

- make a sequence of choices
- at each decision point, pick the choice that seems best at the time

### Greedy vs. DP

- DP also makes a sequence of choice
- current choice depends on earlier choices

$$\text{(ex: } C[i,j] = C[i-1,j] + 1 \text{)}$$

### Risky Pieces

- Pick DP solution when greedy would work
  - ↳ maybe greedy would work too, maybe more efficient (time/space)
- Pick greedy when DP is needed
  - ↳ greedy might not give optimal solution

### (ex) Scheduling problem

greedy choice: earliest finish time

optimizing for: # activities

↳ is this right?

|

↳ reminder :)

$S_k$  — subproblem

$z_m \in S_k \rightsquigarrow$  earliest finish time

there is an optimal solution  $A_k$

• is  $z_m \in A_k$ ? If so, great!

• if not, can we swap  $z_m$  into  $A_k$ ?

$z_j \in A_k$  with earliest finish time

$$f[j] \geq f[m]$$

so, safe to swap in  $z_j, z_m$

But if we picked earliest start time?

$z_m$  has earliest start time

$z_j$  has earliest start time in optimal

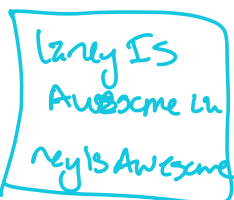
• can we swap  $z_m, z_j$ ? no "h

b/c  $z_m$  might

finish too late

## 2. Huffman Codes

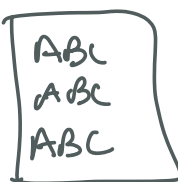
- used to compress data
- In general, we want to rep the same information, but ideally in less space


ex:   
file.txt

Laney Is Awesome  $\times 1000$   
total characters:  $(5+2+7) \times 1000$   
14,000 chars

Compress - replace something with something shorter

Laney - A  
Is - B  
Awesome - C

file 

ABC  $\times 1000$   
total chars: 3000  


File compression as optimization problem

- size of file (smaller is better)
- replace a character (a-z) with something smaller
  - one char stored as one byte (8 bits)
  - goal: use  $\leq 8$  bits for most/some chars  
↳ 0/1

Huffman goal:

- replace char with bits
- based on frequency of characters

• for every char in file, we know its frequency

ex Alphabet

A - 7

B - 11

C - 8

one has 1 bit (0)

two have 2 bits (10)

(11)

B should have 1 bit  $\rightarrow$  most frequent!

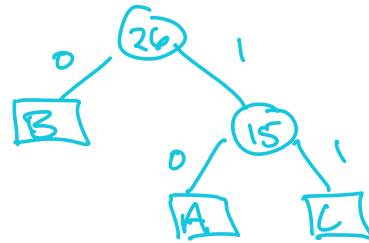
A, C can have 2

$\hookrightarrow$  create a binary tree

characters are leaves

left = 0

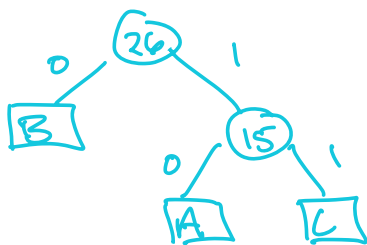
right = 1



$\rightarrow$  code  
 B : 0  
 A : 10  
 C : 11

Prefix-free code:

- no code is a prefix of another code
- 110 / 111 not valid



ex)

11100011  
 C A B B C

(no ambiguity!)

$\rightarrow$  compressed file

10:47

### 3. Greedy Huffman

- given: list of chars with frequencies

- goal: build a tree to rep. an optimal prefix-free code

OS  
Σ

- the tree can be used to build a new, smaller file

greedy choice:

pick the two chars with lowest freqs, build tree from bottom up

### High-level Pseudocode

- start with chars/freqs

- remove the two lowest-freq chars ( $x, y$ )

- merge  $x, y$  together into a subtree

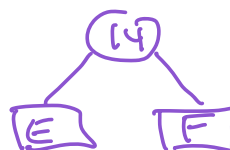
- $x \cdot \text{freq} + y \cdot \text{freq}$  added back into char/freq list

ex Alphabet: a-f

Freqs:      A      B      C      D      E      F  
                 45      13      12      16      9      8

1 bit!  
goze

① Remove 2 lowest-freq chars  
merge into subtree



Put 14 into list

Freq:      A      B      C      D      EF  
                 45      13      12      16      14

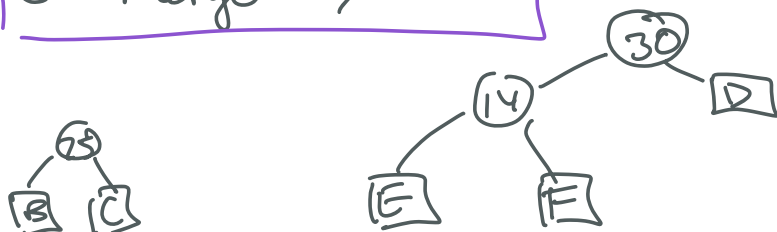
② Remove 2 lowest-freq chars



Put 25 into list

A    BC    D    EF  
45    25    16    14

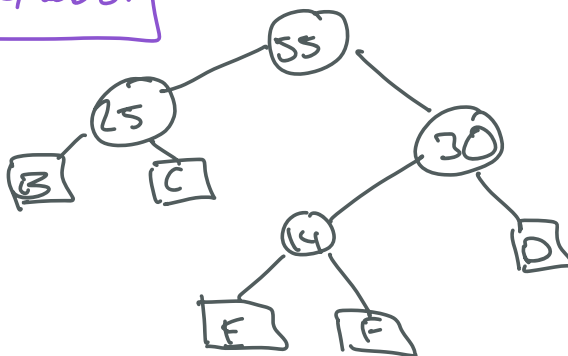
3. merge D, EF



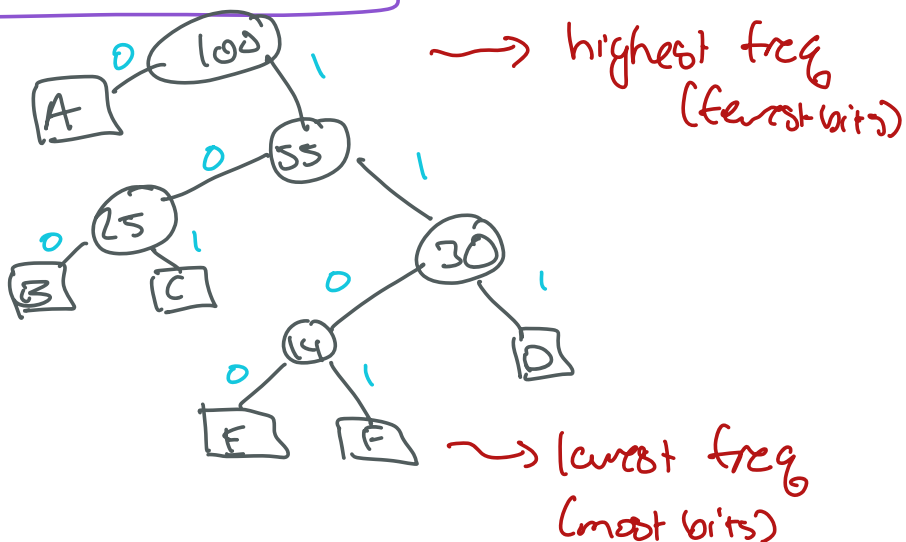
Freqs: A    BC    DEF  
         45    25    30

4. merge BC, & DEF

Freqs: A: 45, BCDEF: 55



5. merge A with rest of tree



ex encode ABC 0100101

ex decode 111100 DB

How much space did we save?

orig: 100 chars, 8 bits each = 800

compressed:  $1 \cdot 45 + 3 \cdot 13 + 3 \cdot 12 + 3 \cdot 16 + 4 \cdot 9 + 4 \cdot 5$   
 A B C D E F

= 224 bits



800 → 224

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