

Mergesort Algorithm

Mergesort is a divide-and-conquer recursive algorithm. It *divides* by making the input array smaller and smaller and smaller until it's trivially easy to sort (sorting the trivially easy is the *conquer* step). Then it merges together the smaller sorted subarrays – that's the final step, which we call *combine*.

MERGESORT(A, p, r)

```

1  if  $p < r$ 
2       $q = \lfloor (p + r) / 2 \rfloor$ 
3      MERGESORT( $A, p, q$ )
4      MERGESORT( $A, q + 1, r$ )
5      MERGE( $A, p, q, r$ )

```

MERGE(A, p, q, r)

```

1   $n_L = q - p + 1$ 
2   $n_R = r - q$ 
3  let  $L[1 : n_L]$  and  $R[1 : n_R]$  be new arrays
4  for  $i = 1$  to  $n_L$ 
5       $L[i] = A[p + i - 1]$ 
6  for  $j = 1$  to  $n_R$ 
7       $R[j] = A[q + j]$ 
8   $i = 1$ 
9   $j = 1$ 
10  $k = p$ 
11 while  $i \leq n_L$  and  $j \leq n_R$ 
12     if  $L[i] \leq R[j]$ 
13          $A[k] = L[i]$ 
14          $i = i + 1$ 
15     else
16          $A[k] = R[j]$ 
17          $j = j + 1$ 
18      $k = k + 1$ 
19 while  $i \leq n_L$ 
20      $A[k] = L[i]$ 
21      $i = i + 1$ 
22      $k = k + 1$ 
23 while  $j \leq n_R$ 
24      $A[k] = R[j]$ 
25      $j = j + 1$ 
26      $k = k + 1$ 

```

Here's how we typeset the above functions in CLRS style:

```

\begin{codebox}
\Procname{$\proc{Mergesort}(A, p, r)$}
\li \If $p < r$
\Then
\li $q \gets \lfloor (p+r)/2 \rfloor$
\li $\proc{Mergesort}(A, p, q)$
\li $\proc{Mergesort}(A, q+1, r)$
\li $\proc{Merge}(A, p, q, r)$
\end{codebox}

```

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```

\begin{codebox}
\Procname{$\proc{Merge}(A, p, q, r)$}
\li $n_L = q - p + 1$
\li $n_R = r - q$
\li let $L[1:n_L]$ and $R[1:n_R]$ be new arrays
\li \For $i \gets 1$ \To $n_L$
\Do
\li $L[i] \gets A[p+i-1]$
\End
\li \For $j \gets 1$ \To $n_R$
\Do
\li $R[j] \gets A[q+j]$
\End
\li $i \gets 1$
\li $j \gets 1$
\li $k \gets p$
\li \While $i \le n_L$ and $j \le n_R$
\Do
\li \If $L[i] \le R[j]$
\Then
\li $A[k] \gets L[i]$
\li $i \gets i + 1$
\li \Else
\li $A[k] \gets R[j]$
\li $j \gets j + 1$
\End
\li $k \gets k + 1$
\End
\li \While $i \le n_L$
\Do
\li $A[k] \gets L[i]$
\li $i \gets i + 1$
\li $k \gets k + 1$
\End
\li \While $j \le n_R$
\Do

```

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\li $A[k] \gets R[j]$  
\li $j \gets j + 1$  
\li $k \gets k + 1$  
\End  
\end{codebox}
```