

CS3000

5/21 - Weds !!

Admin

- exam #1 tmw

- hw2 due 5/22 9pm

- hw3 at tmw, due 5/30 9pm

Agenda

1. linked list

2. Hash table

3. Collisions / Run-Time

4. Exam Prep

1. Linked Lists

Can we do better?

1. O+C

2. data structure

3. ...

trees (pointers)

heap

stacks

queues

(underlying array)

plain array A, length n

starting point
for run-time

w.c.

find

unsorted

$\theta(n)$

sorted

$\theta(\lg n)$

insert

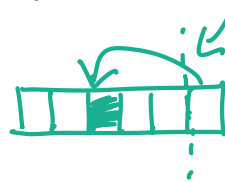
$\theta(1)$

$\theta(n)$

remove

$\theta(1), \theta(n)$

$\theta(n)$



linked list:

- array like
- no positions
- traverse using pointers

head



(assume unsorted)

Each node n:

- n.key
- n.next (could be nil)
- n.prev (could be nil)

Find — basically, linear search

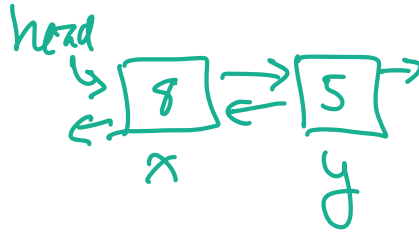
- no positions
- not sorted

- visit every node in order
- compare n.key to key

$\Theta(n)$

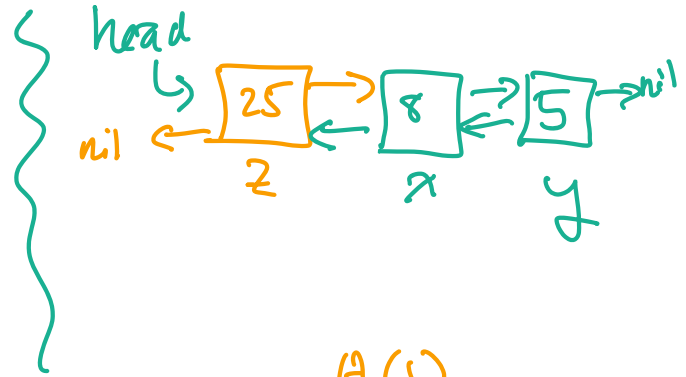
insert

new node z



updates:-

- head points to z
- z.prev is nil
- z.next is x
- x.prev is z

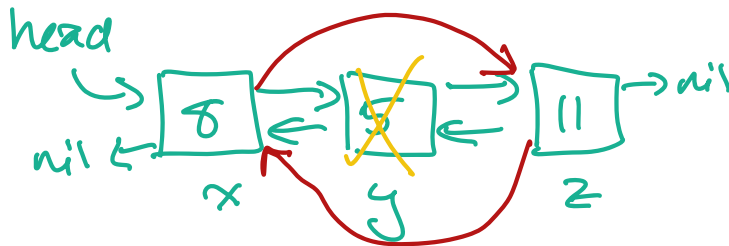


$\Theta(1)$

(and can grow arbitrarily)

remove, once found

- used find to retrieve node to remove
- or, given the node



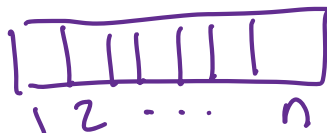
ex) remove y
x.next to y.next
z.prev to y.prev
(deallocate y)

$\Theta(1)$

+ maintains original order!

2. Hash Table

- ↳ underlying array
 - array uses
 - positions



position is determined by the key,
not the order in which we insert

key position

!

2

♥

5

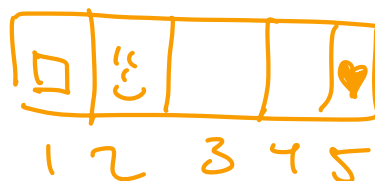
□

1

insert (♥)

insert (□)

insert (!)



Can this be faster
than a regular array?

* do keys have to
be distinct?

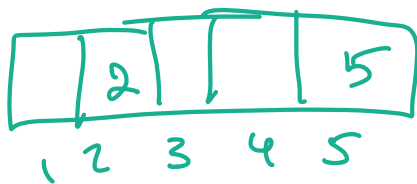
Simplifying assumption: keys are ints ≥ 1

- map key to position (hash function)
 - ↳ key = position?

insert(5)

insert(2)

insert(521,000)



↳ exceeds capacity of array
or, little used capacity

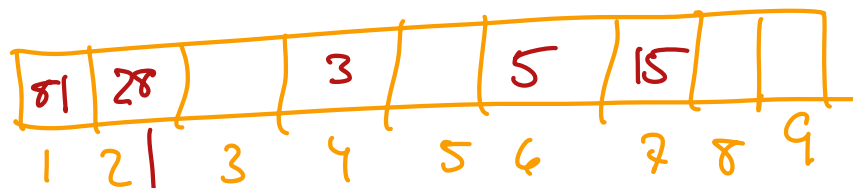
~~521,000~~
521,800

Problem to solve #1: key $\rightarrow$ position gets positions that are too big!

- still have ints ≥ 1 keys
- restrict range of possible values: mod

Ex) hash function $h(k) = k \bmod 9 + 1$

insert:	pos:
5	6
28	2
81	1
3	4
15	7
19	2



something already there!

By introducing mod

- solve problem #1
- now, problem #2! $\vdots$

Collision

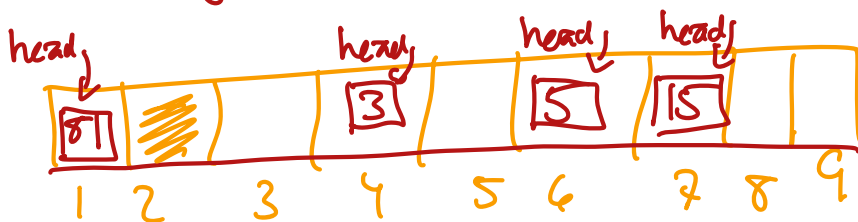
3. Collision + Run-Time

↳ hash function should minimize collisions, but can't avoid them!

• Has to put multiple things at same position in an array?

Linked List at each position in array

- "chaining"
- element first to its position, create a one-element linked list
- in case of collision, add to existing linked list at that position



Assumptions:

- keys don't need to be pos. ints
- hash function is $\Theta(1)$

Insert

- hash key $\Theta(1)$
- create/add to linked list $\Theta(1)$

Find

- hash key $\Theta(1)$
- find in linked list $\Theta(n)$

~~~~~> worst case: all  $n$  elements hashed to the same position !!

## remove

- hash key  $\Theta(1)$
- remove it from the linked list  
↳ once found:  $\Theta(1)$

- not luck
- Bad hash function!

## What makes a good hash function?

- don't put everything in same position
- restricts the range
- predictable/deterministic
- don't waste space / use all available positions
- all positions are equally utilized
- work on different data types

↓  $\Theta(m)$  where  $m$  is length of linked list

Real life: w/ a good hash function!

expected  $\Theta(1)$  for find !!

Hash Table IRL:    / / /  
 $\Theta(1)$  insert, remove, find  $\Xi$   
//       //       -

# 4. Exam #1

## Logistics

- 11:40-1:20 5/22 (be on time!)
- take your time!
- don't write on backs of pages
- pseudocode on there for specific questions
- 8.5x11 inch cheat sheet, one side only
- no phones, calcs, devices, notes etc.
- time displayed on screen

## Questions

- prove correctness
  - iterative run-time
  - recursive run-time
- } given pseudocode,
- compute  $T(n)$
  - (solve  $T(n)$ )
  - give bound
- given  $T(n)$ ,
- solve ~~at~~ 2nd
  - give bound
- no master method
  - use iteration method to solve recurrence
  - put recursive sequence in closed form
  - given pseudocode, what does it do?
  - no space complexity !!
  - write pseudocode, would be straightforward
  - growth of functions ( $O$ ,  $\Omega$ )
  - what does this data structure look like?  
stacks, queues, trees - simple

arrays, heaps — more involved

- given an algorithm we know, what does it do?
- High level: linear search      insertion sort      quicksort  
                                binary search      mergesort      heapsort