CS3000: Algorithms & Data — Summer 2025 — Laney Strange

Homework 1 Due Thursday May 15 @ 9pm via Gradescope

Name: Collaborators:

- Put your name on the first page. If you are using the LATEX template we provided, then you can make sure it appears by filling in the yourname command.
- This assignment is due Thursday May 15 @ 9pm via Gradescope. You may submit up to 48 hours late for no penalty, but expect a delay in grading.
- Show ALL your work, even if the problem doesn't specify it.
- You'll have an opportunity to resubmit one homework at the end of the semester.
- Solutions must be typeset, preferably in LATEX. If you need to draw any diagrams, you may draw them by hand as long as they are embedded in the PDF. I recommend using the source file for this assignment to get started.
- Any solutions that include pseudocode must be in the CLRS style.
- I encourage you to work with your classmates on the homework problems. *If you do collaborate, you must write all solutions by yourself, in your own words.* Do not submit anything you cannot explain. Please list all your collaborators in your solution for each problem by filling in the yourcollaborators command.
- If you get stuck on a homework problem, come by office hours or post on Piazza! We recommend you spend about 30 minutes trying to figure out a problem, and then ask for help. We'll be happy to clarify material from class and algorithm concepts, but we will not give out solutions or confirm your answers are correct.
- Finding solutions to homework problems online, or by speaking with students not enrolled in the class, is strictly forbidden.

Problem 1. Run-Time (5 points)

(a) Express the function $n^3/1000 + 100n^2 - 100n + 3$ in terms of Θ notation.

Solution:

- (b) List the following functions in increasing order of growth (slowest-growing to fastest-growing). As usual, lg means log₂:
 - (a) 5^{5^n}
 - (b) 5*n*
 - (c) 5^{5n}
 - (d) 41g*n*
 - (e) 21g1g*n*
 - (f) *n*⁶
 - (g) $5(n \lg n)$
 - (h) 5^{*n*}
 - (i) 5^{n^5}
 - (j) $(n/6)^6$

Problem 2. Mystery Algorithm (10 points)

You encounter the following pseudocode, and you can assume it calls the procedure LINEARSEARCH implemented as we covered in class. The algorithm accepts an array of n distinct integers in the range 1..n + 1.

MYSTERY(A, n)

- for *i* = 1 to *n* + 1
 if LINEARSEARCH(*A*, *n*, *i*) == NIL
 return *i*
 - (a) In English, describe what this algorithm is trying to accomplish.

Solution:

(b) Formally show the worst-case run-time T(n) of the Mystery Algorithm. Assume that each execution of the *k*th line takes c_k time where c_k is a constant, and assume that LINEARSEARCH itself takes *n* steps.

Start by filling in the two rightmost columns of the table below.

line	cost	times
1		
2		
3		

Solution:

(c) Compute T(n) based on your table above.

Solution:

(d) Give a tight bound on the run-time of this algorithm in the worst-case.

(e) Give an example of *A* that would result in the best-case run time. (Your *A* example must be non-empty.) Give a bound on that run-time.

Problem 3. Improving Run-Time (10 points)

Still assuming that the mystery algorithm above accepts an array of distinct integers in the range 1..n + 1.

(a) Give complete pseudocode, in the CLRS style, for an algorithm that will always return the same result, but with a more efficient run-time in the worst case. (But, sorting the array first is not an acceptable solution!).

Solution:

(b) Compute the worst-case run-time T(n) of your algorithm, showing all your steps.

Solution:

(c) Give a tight bound on the worst case run-time.

Problem 4. Algorithmic Correctness (10 points)

It's here again... our algorithm for LINEARSEARCH.

```
LINEARSEARCH(A, n, key)
```

```
1 for i = 1 to n
2 if A[i] == key
3 return i
```

```
4 return NIL
```

Prove that LINEARSEARCH is correct by showing that the following loop invariant: At the start of each iteration of the for loop, the subarray A[1..i - 1] consists of the elements in A that are not equal to the key. Use the loop invariant proof structure of (1) initialization, (2) maintenance, and (3) termination.

Problem 5. Recursive Sequences (10 points)

- (a) Find the four terms $a_2, ..., a_5$ of the sequence defined by the following recurrence relations.
 - $a_n = 6a_{n-1}$. Base case: $a_1 = 2$ Solution:
 - $a_n = a_{n-1}^2$. Base case: $a_1 = 2$. Solution:
 - $a_n = a_{n-1} + 3a_{n-2}$. Base cases: $a_0 = 1, a_1 = 2$.



(b) Use the iteration method to construct a closed-form formula for an arbitrary term a_k for the first recursive sequence above.

Solution:

(c) Give a recursive formula, including the base case, for the closed-form formula $a_k = \frac{(k)(k+1)}{2}$. Solution: