## CS3000: Algorithms & Data — Summer 2025 — Laney Strange

## APP 7

Due: June 5th, 2025 @ 11:30am via Gradescope

Name: Sample Solution

- APPs will be assigned towards the end of roughly two lectures each week. You'll put together a solution to a short problem that we'll all use in the following lecture. We'll have time set aside to do these in class, or you can work on your own.
- You may handwrite your solutions, or typeset them in LATEXor another system.
- APPs will be graded on completeness. They must be submitted by 11:30am (just before lecture) on the due date. They will not be accepted late, but we drop 3 of them (out of 8 total).
- Collaboration is strongly encouraged for APPs!

## Problem 1. DFS

Topological sort only works on a directed, acyclic graph. A *back edge* is an edge (u, v) connecting a vertex u to an ancestor v in a depth-first tree. Self-loops are considered back-edges. We want to prove that:

- if a directed graph *G* is acyclic, then a DFS of *G* yields no back edges.
- if a DFS of a directed graph G yields no back edges then G is acyclic, and

The first part of the proof is below. Your job is to add the second part.

(Part 1) We will show that if a directed graph *G* is acyclic, then a DFS of *G* yields no back edges. Proof by contrapositive: show that if a DFS of *G* produces a back edge, then *G* has a cycle.

Suppose a DFS search produces a back edge (u, v). Then vertex v is an ancestor of u in the DFS forest. Thus G contains a path from v to u and the back edge (u, v) would complete a cycle.

## **Solution:**

(Part 2) Show that if a DFS of a directed graph *G* yields no back edges, then *G* is acyclic. Proof by contrapositive: show that if *G* has a cycle, then a DFS on *G* yields at least one back edge.

Suppose G contains a cycle c. Let v be the first vertex to be discovered in c, and let (u,v) be the preceding edge in c. At time v.s, the vertices of c form a path of white vertices from v to u. By the white-path theorem, vertex u becomes a descendent of v in the DFS forest and therefore (u,v) is a back edge.