Property-Based Testing

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- Finish data definitions
- Controlling ACL2s
- Property-based testing
- ▶ test? & thm

Data Definitions

- ▶ Demo
- Singleton types
- Recognizers
- Enumerated Types
- Range Types
- Product Types
- Records
- Constructors & Accessors
- Listof Combinator
- Union Types
- Recursive Types
- Data-driven Function Definitions
- Mutually Recursive Data Types

Controlling ACL2s

- :program mode turns off theorem proving in ACL2s
 - no termination analysis is attempted
 - ACL2s will still test contracts and report any errors it finds
 - useful for prototyping & experimenting
- Iogic mode is the default mode and allows you to switch back
 - you cannot define :logic mode functions if they depend on :program mode functions
- Other useful settings (we used some of them in HWK 2)
 - > (acl2s-defaults :set testing-enabled nil)
 - > (set-defunc-termination-strictp nil)
 - > (set-defunc-function-contract-strictp nil)
 - > (set-defunc-body-contracts-strictp nil)
- Documentation via : doc

Property-Based Testing

(definec even-natp (x :nat) :bool
(natp (/ x 2)))

(definec even-intp (x: int) :bool (integerp (/ x 2)))

Here is how we test properties in ACL2s

(test? (implies (natp n)
 (equal (even-intp n)
 (even-natp n))))

This is a property

This gives us way more power than check= as ACL2s checks the property on a large number of examples (user-controlled); passes iff all checks pass

Theorem Proving

(definec even-natp (x :nat) :bool
(natp (/ x 2)))

(definec even-intp (x :int) :bool
(intp (/ x 2)))

Here is a way of proving theorems in ACL2s (thm (implies (natp n) (equal (even-intp n) (even-natp n))))

This gives us way more power than check= & test? as it corresponds to an *infinite* number of checks, i.e., it is always true; passes if ACL2s proves it

Thm vs Test? on Properties



test? vs thm

▶ test?

- user only has to provide properties
- ACL2s does the rest: it automatically looks for counterexamples
- a kind of "light-weight" formal methods
- ▶ thm
 - highest assurance level
 - guarantees that the property is always true
 - requires human expert to interactively drive theorem prover
 - a kind of "heavy-weight" formal methods
- Practical considerations
 - start with test?'s and convert to thm's based on budget/risk analysis

Structure of Properties

- Structure of test?/thm is (test?/thm (implies H C))
- Where H (Hypothesis) is of the form (and (R1 x1)...(Rn xn)...)
 - ▶ All the Ris are recognizers & the xis are variables in C (Conclusion)
 - The second ... can be some other, extra assumptions
- C is a boolean expression
- ▶ We must perform contract checking on all the non-recognizers in H
 - The stuff after the recognizers must satisfy its contracts, assuming everything before it holds
- We must perform contract checking for C
 - All functions in C must satisfy their contracts assuming H holds



(definec even-natp (x :nat) :bool (definec even-intp (x :int) :bool (natp (/ x 2)))

(intp (/ x 2)))

(test? (implies (natp n) (equal (even-intp n) (even-natp n)))) Contract checking passes

(test? (implies (intp n) (equal (even-intp n) (even-natp n)))) Contract checking fails (even-natp n) requires n to be a nat



(definec even-natp (x :nat) :bool (natp (/ x 2)))

(definec even-intp (x :int) :bool (intp (/ x 2)))

(test? (implies
 (and (natp n) (< 20/3 n))
 (equal (even-intp n)
 (even-natp n))))</pre>

Contract checking passes < knows n is a rational

(test? (implies (< 20/3 n) (equal (even-intp n) (even-natp n)))) Contract checking fails < does not know that n is a rational



(definec even-natp (x :nat) :bool (natp (/ x 2)))

(definec even-intp (x :int) :bool (intp (/ x 2)))

Does this property hold?

(definec even-intp (x :int) :bool (if (natp x) (even-natp x) (even-natp (* x -1))))

The two properties characterize even-intp in terms of even-natp, so they show another way we could have defined even-intp

Unifying Observation

(definec even-natp (x :nat) :bool
(natp (/ x 2)))

(definec even-intp (x :int) :bool (intp (/ x 2)))

Contract checking a test?/thm is equivalent to contract checking functions For example, contract checking the test? is equivalent to checking the function

(test?	(implies	(defunc test1 (n)
	(and (natp n) (< 20/3 n))	<pre>:input-contract (and (natp n) (< 20/3 n))</pre>
	(equal (even-intp n)	:output-contract (booleanp (test1 (n))
	(even-natp n)))	(equal (even-intp n)
		(even-natp n))))

In ACL2s, the specification language and programming language are the same!



- Examples in the slides
- More examples

Next Time

- Property-Based Testing in Industry
- Fuzzing for Security Testing
- Next Week: Propositional Logic
- Read Chapter 3