

SOLUTIONS

CS1800

Fall 2025

Recitation 3 - Practice Questions for Homework 2

September 24 & 25, 2025

Recitations

CS1802 Recitations are dedicated time set aside to work on practice problems that specifically prepare you for the current homework or upcoming quiz.

Recitations are in-person and attendance is expected.

The solutions are published at the same time as the problems, so you can check your work. There is no need to submit anything.

Approaching the Problems

These practice problems are labelled according to which Homework or Quiz topic they will help you prepare for. You do not need to complete every practice question; we encourage you to do at least one per topic, and to prioritize the topics you would like to practice.

Instructors & Teaching Assistants

Your recitation is led by a Khoury College professor, assisted by a knowledgeable and wonderful Teaching Assistant. Professors and TAs are fantastic resources, and you have the opportunity in recitation to work with them in a smaller group -- I strongly recommend you take advantage of the time to review your solutions to these practice problems, ask for help on the homework, or review material from lecture.

SOLUTIONS

Practice Problems for Unsigned Numbers (HW2, Question 1)

Convert each of the following (unsigned) numbers from its original base to the given base.

A $010001_2 = \underline{\hspace{2cm}}_{10}$

Solution

$$\begin{aligned} & 0 \cdot 2^5 + 1 \cdot 2^4 + 0 \cdot 2^3 + 0 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0 \\ &= 0 + 16 + 0 + 0 + 0 + 1 \\ &= 17_{10} \end{aligned}$$

B $418_{10} = \underline{\hspace{2cm}}_2$

solution

$$\begin{aligned} 418 \div 2 &= 209 R 0 \\ 209 \div 2 &= 104 R 1 \\ 104 \div 2 &= 52 R 0 \\ 52 \div 2 &= 26 R 0 \\ 26 \div 2 &= 13 R 0 \\ 13 \div 2 &= 6 R 1 \\ 6 \div 2 &= 3 R 0 \\ 3 \div 2 &= 1 R 1 \\ 1 \div 2 &= 0 R 1 \end{aligned}$$

Remainders from bottom to top: 110100010

$$\begin{aligned} \text{Sanity-check to confirm: } 110100010 &= 1 \cdot 2^8 + 1 \cdot 2^7 + 1 \cdot 2^5 + 1 \cdot 2^1 \\ &= 256 + 128 + 32 + 2 = 418 \text{ woohoo :)} \end{aligned}$$

C $42011_5 = \underline{\hspace{2cm}}_{10}$

Solution

$$\begin{aligned} & 1 \cdot 5^0 + 1 \cdot 5^1 + 0 \cdot 5^2 + 2 \cdot 5^3 + 4 \cdot 5^4 \\ &= 1 + 5 + 0 + 250 + 2500 \\ &= 2756_{10} \end{aligned}$$

D $10011100_2 = \underline{\hspace{2cm}}_{16}$

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Solution

Converting from base 2 to base 16, we can use the lookup table. Break into chunks of 4

1 0 0 1 | 1 1 0 0

$$1001 = 9_{16}$$

$$1100 = C_{16}$$

Put them together to get: $9C_{16}$

E $3AB2_{16} = \underline{\hspace{2cm}}_2$

Solution

$$3_{16} = 0011_2$$

$$A_{16} = 1010_2$$

$$B_{16} = 1011_2$$

$$2_{16} = 0010_2$$

String them together to get: 0011101010110010

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Practice Problems for Two's Complement (HW2, Question 2)

- A** What are the highest positive and lowest negative values we can store in 5-bit two's complement? Give your answer in both two's complement binary and decimal.

Solution

- Highest positive: $01111_2 = 15_{10}$
- Lowest Negative: $10000_2 = -16_{10}$

- B** Convert both decimal numbers to 5-bit two's complement, perform the arithmetic in the given base (output must have the same number of digits as input), and report the final answer in decimal.

$$13_{10} + 2_{10}$$

Solution (expanding everything to 5 bits)

$$13_{10} = 01101$$

$$2_{10} = 00010$$

$$\text{Add: } 01111$$

Sanity-check: $01111_2 = 15_{10}$ as we just saw above!

- C** Convert both decimal numbers to 5-bit two's complement, perform the arithmetic indicated, and report the final answer in decimal. If overflow occurs, state it in your answer. (Recall that we don't actually perform the subtraction operation; we add negative values instead!)

$$13_{10} - 10_{10}$$

Solution

$$13_{10} = 01101$$

$$10_{10} = 01010$$

But we want the 10 to be negative, so complete the conversion:

- Flip the bits: 10101
- Add one to get -10 in two's complement: 10110

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$$\begin{array}{r}
 \text{Add} \quad 01101 \quad (13) \\
 + 10110 \quad (10) \\
 \hline
 100011
 \end{array}$$

Oh wait, but that's six bits! Is it overflow?

Nope, we ALWAYS chop off the extra bit. In this case, we chop it off and get the correct answer, so everything is cool.

$$00011_2 = 3_{10} \quad \text{woohoo! :)}$$

NOT OVERFLOW :)

- D** Convert both decimal numbers to 5-bit two's complement, perform the arithmetic indicated, and report the final answer. If overflow occurs, state it in your answer.

$$13_{10} - 2_{10}$$

Solution

$$\begin{array}{r}
 13_{10} = 01101_2 \\
 2_{10} = 00010_2
 \end{array}$$

But we want 2 to be negative, so we can perform the calculation as $13 + -2$.

Flip the bits: 11101

Add one to get -2 in two's complement: 11110

$$\begin{array}{r}
 \text{Add} \quad 01101 \quad (13) \\
 + 11110 \quad (-2) \\
 \hline
 101011
 \end{array}$$

Oh wait, but that's six bits! Is it a problem?

Nah because we can chop off that extra bit and get the correct answer:

$$01011_2 = 11_{10} \quad \text{woohoo! :)}$$

NOT OVERFLOW :)

- E** Would the following operation result in overflow in 5-bit two's complement?

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$$01110 + 00110$$

Yes.

Doing the addition directly gives us 10100, which is clearly a problem, because we've added two positive numbers and gotten a negative (we know its negative because that leading bit is 1 in the output). Must be overflow!

Another explanation: If you convert both to decimal you'll see that the result $14+6=20$ is too big to fit in 5 bits in the two's complement world.

F Would the following operation result in overflow in **unsigned** binary?

$$01110 + 00110$$

Nope! Doing the addition directly gives us 10100_2 , which is 20_{10} --- so as long as we're in an **unsigned** universe it works and still fits in 5 bits.

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Practice Problems for Base Conversions (HW2, Question 3)

Suppose you are given an **unsigned** number; the base isn't specified, but you know that the base ranges are 2-16 (i.e., the smallest base you'll see is 2 and the largest is 16). These unsigned numbers might use the digits 0-9 and the characters A-F.

- A** Is it possible that the digit 3 appears in unsigned binary? Why or why not?

Solution - No

The only digits available in unsigned binary are 0 and 1.

- B** What is the **smallest** possible base for the unsigned number 1054?

Solution - Base 6.

The highest digit we can see in that unsigned number is 5, meaning we must be able to represent the values 0-5 in whatever base this is.

- C** What is the **largest** possible base for the unsigned number 10110?

Solution - Base 16.

The values here are just 0s and 1s, so it **could** be a binary number. But 10110 is a perfectly valid number in base 3, base 4, base 5, etc., all the way up to our highest possibility, base 16.

- D** What is the **smallest** possible base for the unsigned number 937A1028?

Solution - Base 11.

The highest character we can see in that unsigned number is A, meaning we must be able to represent the values 0 through 11.

- E** What is the **only** possible base for the following unsigned numbers (both are the same base) if we know that $10 + 11 = 101$?

Solution - Base 2.

We **MUST** be in base 2 at this point, because any higher base would have $1+1=2$ and here we have $1+1 = 10$. Note that above we limited our possible bases from 2 to 16, so we can't go any lower than base 2.

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Fun fact: bases can also be negative! Base -2 is called *negabinary*.

To convert to a negative base like negabinary, we can use the same division algorithm we used for positive bases. However, when dividing by a negative base, we must choose quotients such that the remainder is non-negative. As an example, we will use the division algorithm to convert 10_{10} to negabinary:

$$\begin{aligned} 10 \div -2 &= -5 \, r \, 0 \\ -5 \div -2 &= 3 \, r \, 1 \\ 3 \div -2 &= -1 \, r \, 1 \\ -1 \div -2 &= 1 \, r \, 1 \\ 1 \div -2 &= 0 \, r \, 1 \end{aligned}$$

Confirmed! From this algorithm, we see that $10_{10} = 1110_{-2}$

A Convert 34_{10} to negabinary

Solution:

$$\begin{aligned} 34 \div -2 &= -17 \, r \, 0 \\ -17 \div -2 &= 9 \, r \, 1 \\ 9 \div -2 &= -4 \, r \, 1 \\ -4 \div -2 &= 2 \, r \, 0 \\ 2 \div -2 &= -1 \, r \, 0 \\ -1 \div -2 &= 1 \, r \, 1 \\ 1 \div -2 &= 0 \, r \, 1 \end{aligned}$$

And so $34_{10} = 1100110_{-2}$

B Convert 17_{10} to negabinary

Solution:

$$\begin{aligned} 17 \div -2 &= -8 \, r \, 1 \\ -8 \div -2 &= 4 \, r \, 0 \\ 4 \div -2 &= -2 \, r \, 0 \\ -2 \div -2 &= 1 \, r \, 0 \\ 1 \div -2 &= 0 \, r \, 1 \end{aligned}$$

And so $17_{10} = 10001_{-2}$

C Convert 00111_{-2} to decimal.

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Solution

$$\begin{aligned}00111_{-2} &= 1 \cdot -2^2 + 1 \cdot -2^1 + 1 \cdot -2^0 \\&= 1 \cdot 4 + 1 \cdot -2 + 1 \cdot 1 \\&= 4 - 2 + 1 \\&= 3_{10}\end{aligned}$$

- D** Can a decimal number have the same binary and negabinary representation? Give an example if Yes, or justify if No.

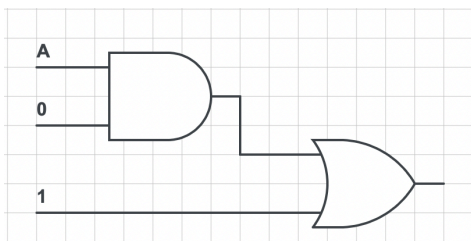
Solution: Yes, for example.... 17, which we just did! $17_{10} = 10001_2 = 10001_{-2}$.

This will be true, not just for 17, but for any value where the binary representation has 1s only on even powers of 2.

Practice Problems for Circuits & Logic (HW2, Question 4)

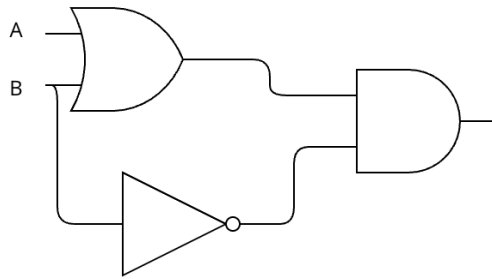
- A** Draw a circuit that represents the logic statement $(A \wedge F) \vee T$. Do not simplify the expression; your circuit should implement the statement directly!

Solution (it's fine to use 0/1 OR F/T in circuits!)



- B** Write the truth table for the circuit shown below, including all intermediate steps.

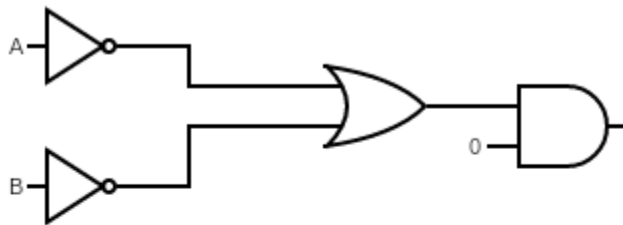
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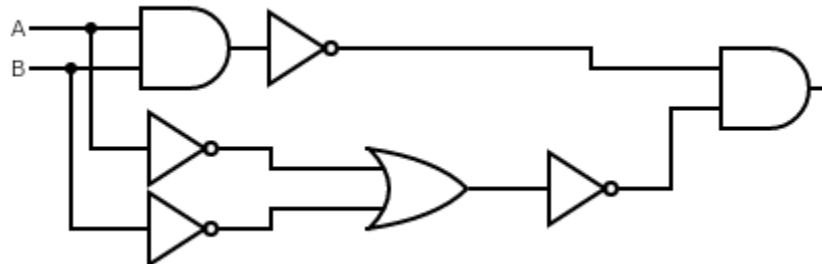
Solution

A	B	$\neg B$	$A \vee B$	$(A \vee B) \wedge \neg B$
0	0	1	0	0
0	1	0	1	0
1	0	1	1	1
1	1	0	1	0

(For Parts C-G, Consider the two circuits below.)



Circuit #1



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Circuit #2

- C** Write out logical expressions for both circuits

Solution

Circuit 1: $(\neg A \vee \neg B) \wedge F$

Circuit 2: $\neg(A \wedge B) \wedge \neg(\neg A \vee \neg B)$

- D** Show that the two circuits are equivalent using the laws of logical equivalence. As always, apply one law at a time, take small steps and label each one.

Proof with logical equivalence laws:

Simplify Circuit 1

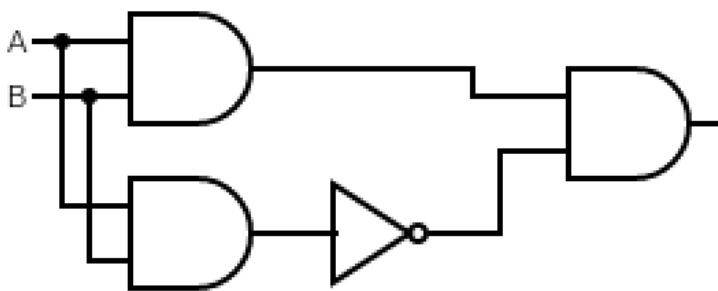
- $(\neg A \vee \neg B) \wedge F$
- F domination

Simplify Circuit 2

- $\neg(A \wedge B) \wedge \neg(\neg A \vee \neg B)$
- $\neg(A \wedge B) \wedge \neg(\neg(A \wedge B))$ deMorgan's complement
- F

- E** There are many other circuits that would be equivalent to these two. Draw one that uses three AND gates, one NOT gate, and no other gates.

Solution



- F** Write out a logical expression for your new circuit.

$(A \wedge B) \wedge \neg(A \wedge B)$.

- G** Simplify your statement by applying the laws of logical equivalence. Clearly state every law applied.

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Note that we start with an expression $(A \wedge B)$ and its negation $\neg(A \wedge B)$, and they are and'ed together to form the circuit above: $(A \wedge B) \wedge \neg(A \wedge B)$. Because of this, we can apply the negation law directly, because an expression and'ed with its negation is always False.

$$(A \wedge B) \wedge \neg(A \wedge B)$$

F

This is the original circuit directly

Negation law

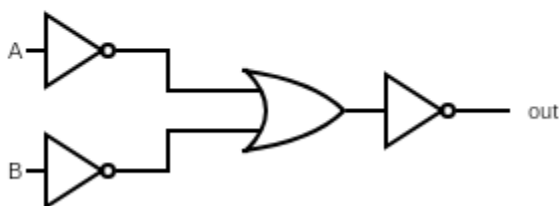
SOLUTIONS

Practice Problems for Circuit Operations (HW2, Question 5)

- A** Sometimes, we might not have the particular logic gate we need available. Luckily, we can replicate the behavior of some logic gates using other gates. Draw a circuit with the same behavior as an AND gate with only the OR and NOT gates.

Write a logical expression that represents your circuit, and use a truth table to confirm it produces the same output as an AND gate.

Sample Solution



In Logic: $\neg(\neg A \vee \neg B)$

A	B	$\neg A$	$\neg B$	$\neg A \vee \neg B$	$\neg(\neg A \vee \neg B)$
0	0	1	1	1	0
0	1	1	0	1	0
1	0	0	1	1	0
1	1	0	0	0	1

(Parts B-E are directly useful for homework 2 and should be completed in order :)

- B** What is the result of adding together the following unsigned bits? Append as many leading bits as necessary to ensure no overflow in this unsigned universe.

$$0 + 0 = \underline{\hspace{2cm}}$$

$$0 + 1 = \underline{\hspace{2cm}}$$

$$1 + 0 = \underline{\hspace{2cm}}$$

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$$1 + 1 = \underline{\hspace{2cm}}$$

Solution:

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

$$1 + 1 = 10$$

- C** Now suppose we have only one bit to store the result of the addition, and if we need more than one bit, we save it as a “carry.”

$$0 + 0 \quad \text{sum } \underline{\hspace{2cm}} \quad \text{carry } \underline{\hspace{2cm}}$$

$$0 + 1 \quad \text{sum } \underline{\hspace{2cm}} \quad \text{carry } \underline{\hspace{2cm}}$$

$$1 + 0 \quad \text{sum } \underline{\hspace{2cm}} \quad \text{carry } \underline{\hspace{2cm}}$$

$$1 + 1 \quad \text{sum } \underline{\hspace{2cm}} \quad \text{carry } \underline{\hspace{2cm}}$$

Solution:

$$0 + 0 \quad \text{sum } \underline{\hspace{1cm}}0\underline{\hspace{1cm}} \quad \text{carry } \underline{\hspace{1cm}}0\underline{\hspace{1cm}}$$

$$0 + 1 \quad \text{sum } \underline{\hspace{1cm}}1\underline{\hspace{1cm}} \quad \text{carry } \underline{\hspace{1cm}}0\underline{\hspace{1cm}}$$

$$1 + 0 \quad \text{sum } \underline{\hspace{1cm}}1\underline{\hspace{1cm}} \quad \text{carry } \underline{\hspace{1cm}}0\underline{\hspace{1cm}}$$

$$1 + 1 \quad \text{sum } \underline{\hspace{1cm}}0\underline{\hspace{1cm}} \quad \text{carry } \underline{\hspace{1cm}}1\underline{\hspace{1cm}}$$

- D** Write those same answers in a truth table, with A and B as input values and *sum* and *carry* as output values. Using only \wedge , \vee , \neg operators, give two logical expressions that represent the truth table (one for sum and one for carry).

A	B	Sum	Carry

Solution

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A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

$$\text{Sum} = (A \wedge \neg B) \vee (\neg A \wedge B)$$

$$\text{Carry} = A \wedge B$$

- E** Great, that's how we do addition with circuits! Now let's expand this so we're adding columns of bits like we did in the binary-number exercises above. Suppose we're adding the two unsigned binary numbers $110 + 111$.

Every column we add has a *carry-in* (from the previous column) and a *carry-out* (to the next column). For each column below, specify the carry-in, sum, and carry-out.

	1	1	0
+	1	1	1

Solution

	1	1	0
+	1	1	1
	Carry-in: 1 Sum: 1 Carry: 1	Carry-in: 0 Sum: 0 Carry: 1	Carry-in: 0 Sum: 1 Carry: 0