

CS1800
11/18 - Tues.

Admin

- HW7 due 12/2 9pm
- Second chance HW due 12/5 9pm (no 48 hour extension)
- Quiz 3 out today
- 11/25 optional lecture
↳ poll on piazza

Agenda

1. Induction Review
2. Inequality ex.
3. Structural ex.

Lecture question

Will we need to be able to do a proof on the quiz?

→ yes! Quiz 4 12/5
one question
30-minute quiz (+ buffer)

1. Induction Review

For induction proofs, how is it considered "proven" if it is based on the ASSUMPTION that the $p(k)$ is true? What about if it wasn't true?

predicate $P(x)$

↳ of time, of interest $P(1), P(2), P(3), P(4) \dots$ (could prove all separately)

induction is a shortcut

↳ $P(k) \Rightarrow P(k+1)$

IF $P(k)$ is true, then $P(k+1)$

↳ IH, assumed

↳ prove, using IH

It needs to be proven true at least once!

↳ Base case!

ex: show $P(1)$

Steps of a proof: ① Predicate ② Logic Statement ③ Base case ④ Inductive step

check (labelled: IH, what is $P(k)$?, goal what is $P(k+1)$?, applying IH

Categories where induction is useful:

1. Summation
2. divisibility
3. inequality
4. structural
5. program correctness

Summation

- def of summation
- formula
- prove: formula always works!

$$\sum_{i=1}^n i = 1 + 2 + \dots + n \quad (\text{def})$$

divisibility

- expression is a multiple of some integer
- if x is a multiple of y , then: $\exists w$ such that $w \cdot y = x$
- If x is a multiple of y , then so is $2x, 3x, \dots$
- If x, z are multiples of y , then xz is a multiple, $x+z$ is a multiple

$$= \frac{(n)(n+1)}{2} \quad (\text{formula})$$

inequality

- some expression $<$ other expression
- ex: $f(n) < g(n)$
- n is input to a function
- induction on n (1, 2, 3, ...)
- we can do anything as long as it's still true! ♥

$$100x < 100y$$

$$100x < 200y$$

$$100x < 200y^2 + 18500$$

structural

- discrete structure property
- set, sequence, graph, etc.
- induction on the size of the structure

revisit to divis. proof from last time
 $6^x + 4$ is a multiple of 5

Predicate $P(x)$ states that $6^x + 4$ is a multiple of 5 $\sim \exists w \in \mathbb{Z}$ such that $5 \cdot w = 6^x + 4$

Logic statement $\forall n \in \mathbb{Z}^+ P(n)$ \sim induction on x

Base case

$P(1)$ $6^1 + 4$ is a multiple of 5
 $6^1 + 4 = 10$ True!



expression: $6^x + 4$
 plug in 1: $6^1 + 4 = 10$
 10 is a multiple of 5 $10 = 5 \cdot 2$

Inductive step $P(k) \Rightarrow P(k+1)$ **Implication**

Inductive Hypothesis $P(k)$, assumed true: $(6^k + 4)$ is a multiple of 5.

Goal: show $P(k+1)$, i.e., $(6^{k+1} + 4)$ is a multiple of 5.

Start with the expression first:

$$\begin{aligned}
 (6^{k+1} + 4) &= 6^k \cdot 6^1 + 4 && // \text{exponent law } a^b a^c = a^{b+c} \\
 &= 6^k \cdot (5 + 1) + 4 && // 6 = 5 + 1. \text{ We do this because 5 is useful!} \\
 &= 6^k \cdot 5 + 6^k \cdot 1 + 4 && // \text{Distribute the } 6^k
 \end{aligned}$$

$6^k \cdot 1 + 4$ is a multiple of 5, by our Inductive Hypothesis.

$6^k \cdot 5$ is multiple by 5, by definition of multiplication

The sum of two multiples of 5 is itself a multiple of 5. Done!

$6^{k+1} + 4$ is multiple of 5 \sim need to prove!

$6 = 5 + 1$ b/c useful!

IH: $6^k + 4$ is multiple of 5! \rightarrow add!

$6^k \cdot 5$ is multiple of 5

$$\begin{aligned} &\hookrightarrow 6^k(5+1) + 4 \\ &= \underbrace{6^k \cdot 5}_{\text{mult. of } 5} + \underbrace{6^k \cdot 1 + 4}_{\text{mult. of } 5 \text{ (IH)}} \end{aligned}$$

2. Inequality Proof

• Predicate $P(x)$ $5x+5 \leq x^2$

• Logic Statement — when is this true?

$$\forall n \in \mathbb{Z}^+ \quad n \geq 6 \Rightarrow P(n)$$

• Base case: prove $P(6)$

$$\begin{aligned} \text{LHS: } &5 \cdot 6 + 5 = 35 && 35 \leq 36 \quad \checkmark \\ \text{RHS: } &6^2 = 36 \end{aligned}$$

$P(1)$ False!

$P(2)$ False!

$P(3)$ False!

$P(4)$ False!

$P(5)$ False!

$P(6)$ True! ;)

↳ and up!

• Inductive Step: $P(k) \Rightarrow P(k+1)$

(IH) $P(k)$, ie, $5k+5 \leq k^2$ — assume

(goal) $P(k+1)$ $5(k+1)+5 \leq (k+1)^2$ — want to prove

Start with LHS:

$$\begin{aligned} 5(k+1) + 5 &= 5k + 5 + 5 && \text{// multiply 5 through} \\ &\leq k^2 + 5 && \text{// by IH} \\ &\leq k^2 + k && \text{// b/c } k \text{ is at least } 6 \\ &\leq k^2 + 2k + 1 && \text{// b/c } k < 2k + 1 \\ &= (k+1)^2 \end{aligned}$$

done!

3. Structural Proof

↳ property of a discrete structure
(set, sequence, graph)

previous set proofs:

• set equality, two sets w/ same elements

induction:

• set property, ignore specific elements

(*) sets!

• induction on cardinality of the set

• $P(1)$ — set with one element

$P(2)$ — set with two elements

...

Focus on: subsets of a set,
of size 2

How many subsets of size 2?

$$P(x) \quad |A|=x, A \text{ has } \frac{(x)(x-1)}{2}$$

subsets of size 2.

$$A = \{a, b\}$$

subsets of size 2: 1

$$A = \{a, b, c\}$$

subsets of size 2: 3

$$A = \{a, b, c, d\}$$

subsets of size 2: 6

Logic Statement: $\forall n \in \mathbb{Z} \quad n \geq 2 \Rightarrow P(n)$

~ could also use $\forall n \in \mathbb{Z}^+$

Base case: $|A|=2 \quad P(2)$

one 2-element subset,
A itself

$$\text{Formula: } \frac{(2)(2-1)}{2} = 1 \quad \checkmark$$

Inductive Step: $P(k) \Rightarrow P(k+1)$

(IH) $P(k)$ is true: $|A|=k$, then A has $\frac{(k)(k-1)}{2}$ 2-element subsets

(Goal) $P(k+1)$ $|A|=k+1$, then A has $\frac{(k+1)k}{2}$ 2-element subsets

• let A be a set with $k+1$ elements, $|A|=k+1$

• remove one element, x , from A, to get set A'

$$A' = A - \{x\} \quad A = \{u, v, \dots, w, x\} \quad |A|=k+1$$

$$A' = \{u, v, \dots, w\} \quad |A'|=k$$

A' has $\frac{(k)(k-1)}{2}$ 2-element subsets (by IH)

• element x is not in any of A' subsets

• create 2-element subsets $\{x, -\}$ ~ How many are there?

↳ total 2-element subsets of A

is 2-element subsets of $A' + \{x, -\}$

↓
pair up x with
every element in A'

↓
there are k of them!

total 2-element subsets in A:

$$A' \text{ 2-element: } \frac{(k)(k-1)}{2} \quad \{x, -\}: k$$

$$\frac{(k)(k-1)}{2} + k$$

$$= \frac{(k)(k-1)}{2} + \frac{2k}{2}$$

$$= \frac{k^2 - k + 2k}{2}$$

$$= \frac{k^2 + k}{2}$$

$$= \frac{(k)(k+1)}{2}$$

done!

// total subsets of size 2

// common term

// terms together

core was:

2-element subsets in A

$$\frac{(k)(k+1)}{2}$$